

# Introduction to Modern Cryptography

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(Slides courtesy of Prof. Jonathan Katz)

Lecture 03, part 1

# Defining Secure Encryption

# Lessons learned so far

## Crypto Design Lesson One

- ▶ The key space must be large enough to make brute-force attacks impractical (cf. Shift Cipher)

## Crypto Design Lesson Two

- ▶ Large key space is a necessary, but not sufficient condition for a secure encryption scheme (cf. Vigenère Cipher)

But what does *secure* actually mean?

## In this lecture

- ▶ What do we mean by **secure**?
- ▶ How do we know when a scheme is **secure**?
- ▶ Can we **prove** that some encryption scheme is **secure**?

# Three principles of modern cryptography

## Definitions

- ▶ Precise, mathematical model and formal definition of what security means

## Assumptions

- ▶ Clearly stated and unambiguous

## Proofs

- ▶ *Prove security* and move away from design-break-patch

# Defining secure encryption

## Security guarantee/goal

- ▶ What we want to achieve (or what we want to prevent the attacker from achieving)

## Threat model

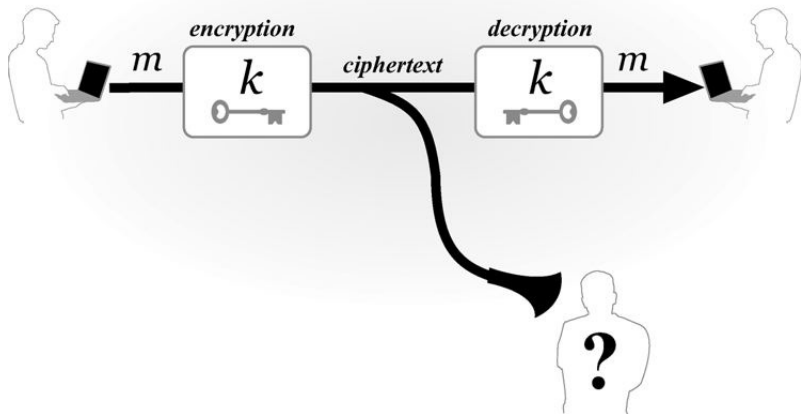
- ▶ What (real-world) capabilities the attacker is assumed to have

## Private-key encryption (recall)

A private-key encryption scheme is defined by a message space  $\mathcal{M}$ , (key space  $\mathcal{K}$ ) and algorithms (Gen, Enc, Dec) :

- ▶ Gen (key-generation algorithm): outputs  $k \in \mathcal{K}$
- ▶ Enc (encryption algorithm): takes key  $k$  and message  $m \in \mathcal{M}$  as input; outputs ciphertext  $c \leftarrow \text{Enc}_k(m)$
- ▶ Dec (decryption algorithm): takes key  $k$  and ciphertext  $c$  as input; outputs  $m$  or "error":  $m = \text{Dec}_k(c)$

# Private-key encryption





# Threat models for encryption

- ▶ Ciphertext-only attack
  - ▶ One ciphertext
  - ▶ Many ciphertexts
- ▶ Known-plaintext attack
- ▶ Chosen-plaintext attack
- ▶ Chosen-ciphertext attack

# Goal of secure encryption?

- ▶ How would you define what it means for encryption scheme  $(\text{Gen}, \text{Enc}, \text{Dec})$  over message space  $\mathcal{M}$  to be **secure**?
- ▶ ...against a (single) ciphertext-only attack

# Secure encryption?

**”Impossible for the attacker to learn the key”**

- ▶ The key is a means to an end, not the end itself
- ▶ Necessary (to some extent) but not sufficient
- ▶ Easy to design an encryption scheme that hides the key completely, but is insecure
- ▶ Can design schemes where most of the key is leaked, but the scheme is still secure

# Secure encryption?

**”Impossible for the attacker to learn the plaintext from the ciphertext”**

- ▶ What if the attacker learns 90% of the plaintext?

# Secure encryption?

**“Impossible for the attacker to learn any character of the plaintext from the ciphertext”**

- ▶ What if the attacker is able to learn (other) partial information about the plaintext?
  - ▶ e.g. salary is greater than **75K**
- ▶ What if the attacker guesses a character correctly?

# The right definition

## Secure encryption

Regardless of any **prior information** the attacker has about the plaintext, the ciphertext should leak **no additional information** about the plaintext

How to formalize?  $\implies$  defining **perfect secrecy** (next slides!)

**End**

Reference: From Section 1.3 until the end of Chapter 1 of the book.