

# **Algorithms and Data Structures**

Upper and Lower Bounds for Sorting, Matrix  
Multiplication

# Matrix Multiplication

Assume that we have two square  $(n \times n)$ -matrices

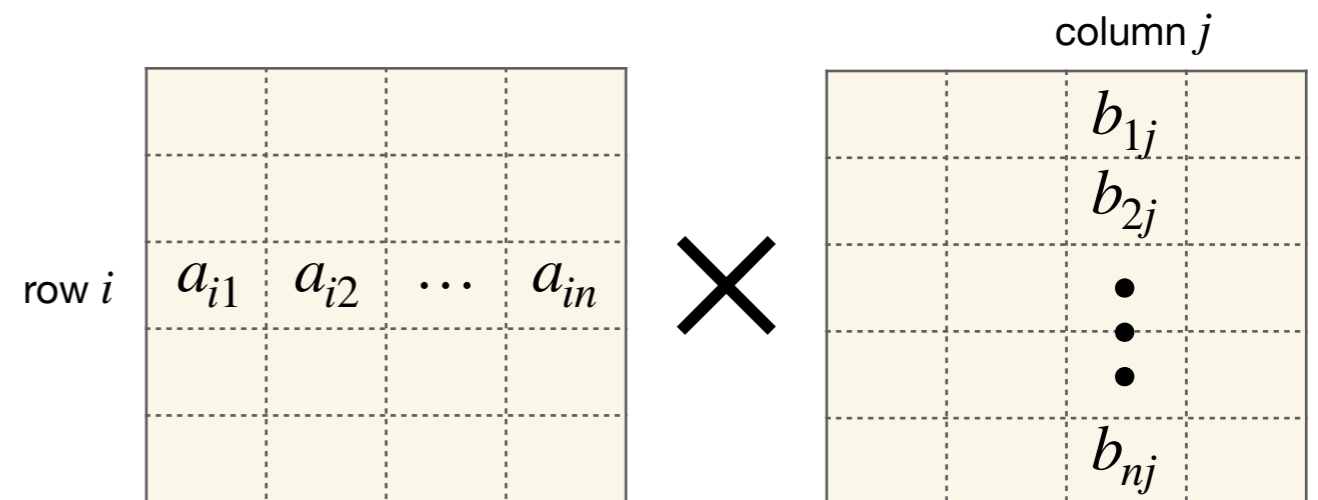
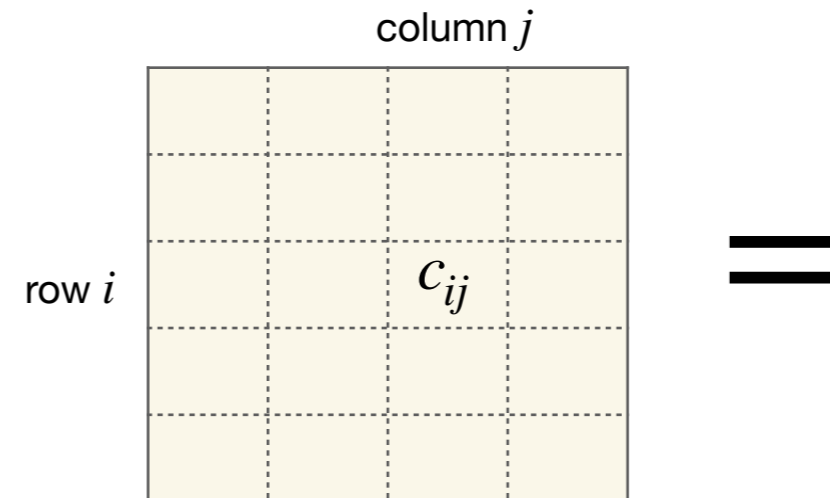
$$A = (a_{ij})_{1 \leq i, j \leq n} \text{ and}$$

$$B = (b_{ij})_{1 \leq i, j \leq n}$$

The product of  $A$  and  $B$  is the  $(n \times n)$ -matrix

$$C = (c_{ij})_{1 \leq i, j \leq n} \text{ with entries}$$

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj}$$



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Can we do better than that?

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For (possibly large matrices), multiplications are more expensive!

# Divide and Conquer...

Suppose we divide our matrices  $A$  and  $B$  as follows:

$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \quad B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

We can write  $C$  as:

$$\begin{aligned} \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} &= \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix} \\ &= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix} \end{aligned}$$

# ...using fewer multiplications.

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

$$P_2 = (A_{21} + A_{22}) \cdot B_{11}$$

$$P_3 = A_{11} \cdot (B_{12} - B_{22})$$

$$P_4 = A_{22} \cdot (-B_{11} + B_{21})$$

$$P_5 = (A_{11} + A_{22}) \cdot B_{22}$$

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$$C_{11} = P_1 + P_4 - P_5 + P_7$$

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$$C_{21} = P_2 + P_4$$

$$C_{22} = P_1 + P_3 - P_2 + P_6$$

# Let's calculate $C_{11}$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

$$P_2 = (A_{21} + A_{22}) \cdot B_{11}$$

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$$P_7 = (A_{12} - A_{22}) \cdot (B_{21} + B_{22})$$

$$C_{11} = P_1 + P_4 - P_5 + P_7$$

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$



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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \times B_{11} + A_{22} \cdot B_{22}$$

$$P_4 = A_{22} \cdot (-B_{11} + B_{21}) = -A_{22} \cdot B_{11} + A_{22} \cdot B_{21}$$

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$$P_1 + P_4 = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21}$$

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$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \times B_{11} + A_{22} \cdot B_{22}$$

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# Let's calculate $C_{11}$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{11} + A_{22} \cdot B_{22}$$

$$P_4 = A_{22} \cdot (-B_{11} + B_{21}) = -A_{22} \cdot B_{11} + A_{22} \cdot B_{21}$$

$$P_5 = (A_{11} + A_{12}) \cdot B_{22} = A_{11} \cdot B_{22} + A_{12} \cdot B_{22}$$

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$$P_1 + P_4 = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21}$$

$$P_1 + P_4 - P_5 = A_{11} \cdot B_{11} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21} - A_{12} \cdot B_{22}$$



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$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{11} + A_{22} \cdot B_{22}$$

$$P_4 = A_{22} \cdot (-B_{11} + B_{21}) = -A_{22} \cdot B_{11} + A_{22} \cdot B_{21}$$

$$P_5 = (A_{11} + A_{12}) \cdot B_{22} = A_{11} \cdot B_{22} + A_{12} \cdot B_{22}$$

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$$P_1 + P_4 = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21}$$

$$P_1 + P_4 - P_5 = A_{11} \cdot B_{11} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21} - A_{12} \cdot B_{22}$$

# Let's calculate $C_{11}$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

$$P_2 = (A_{21} + A_{22}) \cdot B_{11}$$

$$P_3 = A_{11} \cdot (B_{12} - B_{22})$$

$$P_4 = A_{22} \cdot (-B_{11} + B_{21})$$

$$P_5 = (A_{11} + A_{22}) \cdot B_{22}$$

$$P_6 = (-A_{11} + A_{21}) \cdot (B_{11} + B_{12})$$

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$$C_{11} = P_1 + P_4 - P_5 + P_7$$

$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$

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$$P_1 + P_4 = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{22} + A_{22} \cdot B_{21}$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$P_4 = A_{22} \cdot (-B_{11} + B_{21}) = -A_{22} \cdot B_{11} + A_{22} \cdot B_{21}$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$P_5 = (A_{11} + A_{22}) \cdot B_{22}$$

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$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \cdot B_{11} + A_{22} \cdot B_{22}$$

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$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

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$$\begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \times \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$= \begin{pmatrix} A_{11} \cdot B_{11} + A_{12} \cdot B_{21} & A_{11} \cdot B_{12} + A_{12} \cdot B_{22} \\ A_{21} \cdot B_{11} + A_{22} \cdot B_{21} & A_{21} \cdot B_{12} + A_{22} \cdot B_{22} \end{pmatrix}$$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22}) = A_{11} \cdot B_{11} + A_{11} \cdot B_{22} + A_{22} \times B_{11} + A_{22} \cdot B_{22}$$

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Try it at home: Check  $C_{12}$ ,  $C_{21}$ , and  $C_{22}$ .

# Let's calculate $C_{11}$

$$P_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

$$P_2 = (A_{21} + A_{22}) \cdot B_{11}$$

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# How many multiplications do we need?

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We have 7 multiplications  
of  $n/2 \times n/2$  matrices.

$$C_{11} = P_1 + P_4 - P_5 + P_7$$

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Suppose  $T(n) \leq \alpha T(\lceil n/b \rceil) + O(n^d)$

for some constants  $\alpha > 0$ ,  $b > 1$  and  $d \geq 0$ .

$$\text{Then, } T(n) = \begin{cases} O(n^d), & \text{if } d > \log_b \alpha \\ O(n^d \log_b n), & \text{if } d = \log_b \alpha \\ O(n^{\log_b \alpha}), & \text{if } d < \log_b \alpha \end{cases}$$

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Timeline of matrix multiplication exponent

Year	Bound on omega	Authors
1969	2.8074	Strassen <sup>[1]</sup>
1978	2.796	Pan <sup>[10]</sup>
1979	2.780	Bini, Capovani [it], Romani <sup>[11]</sup>
1981	2.522	Schönhage <sup>[12]</sup>
1981	2.517	Romani <sup>[13]</sup>
1981	2.496	Coppersmith, Winograd <sup>[14]</sup>
1986	2.479	Strassen <sup>[15]</sup>
1990	2.3755	Coppersmith, Winograd <sup>[16]</sup>
2010	2.3737	Stothers <sup>[17]</sup>
2012	2.3729	Williams <sup>[18][19]</sup>
2014	2.3728639	Le Gall <sup>[20]</sup>
2020	2.3728596	Alman, Williams <sup>[21][22]</sup>
2022	2.371866	Duan, Wu, Zhou <sup>[23]</sup>
2024	2.371552	Williams, Xu, Xu, and Zhou <sup>[2]</sup>

**Selection**

# The selection problem

**Definition:** The  $i^{\text{th}}$ -order statistic of a set of  $n$  (distinct) elements is the  $i^{\text{th}}$  smallest element.

i.e., the element which is larger than exactly  $i - 1$  other elements.

## The Selection Problem:

**Selection**( $\mathbf{A}[1, \dots, n]$ ,  $i$ )

**Input:** A set of  $n$  (distinct) numbers (in an array  $\mathbf{A}$ ) and a number  $i$ , with  $1 \leq i \leq n$ .

**Output:** The  $i^{\text{th}}$ -order statistic of the set.

**An easy solution**



# An easy solution

Sort the numbers in  $O(n \log n)$  time using MergeSort.

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Sort the numbers in  $O(n \log n)$  time using MergeSort.

Return the  $i^{\text{th}}$  element of the sorted array.

# An easy solution

Sort the numbers in  $O(n \log n)$  time using MergeSort.

Return the  $i^{\text{th}}$  element of the sorted array.

Is sorting an overkill?

# Divide and conquer

Split the input into smaller inputs.

Solve the problem for the smaller inputs recursively.

Combine the solutions into a solution for the original problem.

# The Partition procedure

Procedure **Partition**( $A[i, \dots, j]$ )

Choose a **pivot element**  $x$  of  $A$

$k = i$

For  $h = i$  to  $j$  do

If  $A[h] < x$

Swap  $A[k]$  with  $A[h]$

$k = k + 1$

Swap  $A[k]$  with  $A[h]$

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Running time  **$O(n)$**



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# The Partition procedure (with the pivot element as input)

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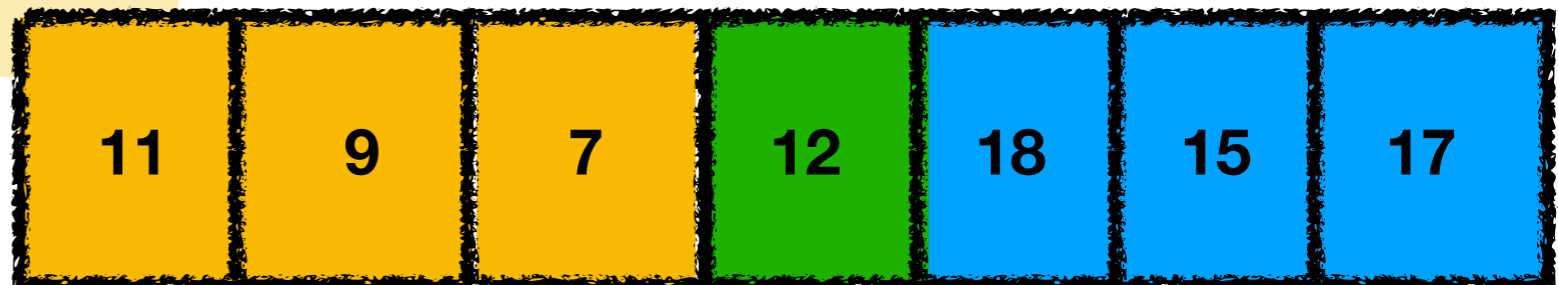
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**Selection**( $\mathbf{A}[1, \dots, n], (n + 1)/2$ )

# Let's try to do that...

Algorithm **Selection**(**A**[1, ...,  $n$ ],  $i$ )

$x = \text{Selection}(\mathbf{A}[1, \dots, n], (n + 1)/2)$

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Do you see a problem?

*x* = **Selection**(**A**[1,..., *n*], (*n* + 1)/2)

*k* = **Partition**(**A**[1,..., *n*], *x*)

Before we conquer, we need to divide!

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We need to partition the array into two using a good pivot element (*the median*).

Or otherwise the running time of the recursion will be bad!

But to find the median, we need an algorithm for selection!

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We need to partition the array into two using a good pivot element (*something “close” to the median*).

Or otherwise the running time of the recursion will be bad!

But to find the median, we need an algorithm for selection!

# A good pivot element

Split the array **A** into sub-arrays with **5 elements** each.

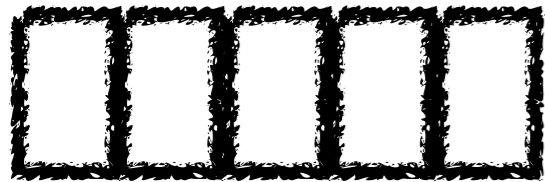
The last one might have fewer elements.

# A good pivot element





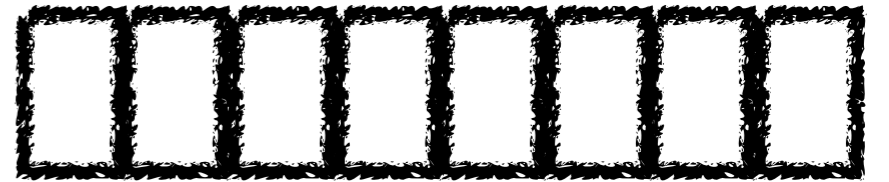
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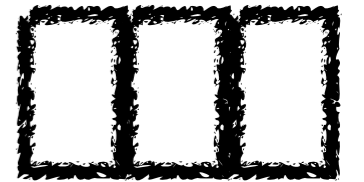
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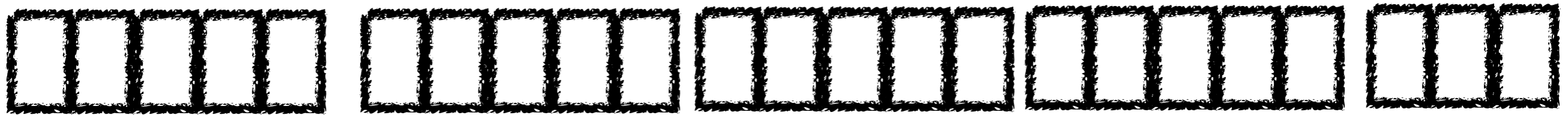
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Run **InsertionSort**

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Find the **median-of-medians**.

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Run **Selection**

# This failed...

Algorithm **Selection**(**A**[1, ...,  $n$ ],  $i$ )

$x = \text{Selection}(\mathbf{A}[1, \dots, n], (n + 1)/2)$

$k = \text{Partition}(\mathbf{A}[1, \dots, n], x)$



# ...but this won't.

Algorithm **Selection**( $\mathbf{A}[1, \dots, n]$ ,  $i$ )

Split the array  $\mathbf{A}$  into  $n/5$  arrays of size 5

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$x = \mathbf{Selection}(\mathbf{A}[m_1, \dots, m_{n/5}], (n/5 + 1)/2)$

/\*Find the median of medians \*/

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# The Selection algorithm

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$k - 1$  is the number of elements in the lower subarray.

If  $i = k$ , return  $x$

If  $i < k$ , return **Selection**( $\mathbf{A}[1, \dots, k - 1]$ ,  $i$ )

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We are looking for the  $i^{\text{th}}$ -order statistic.

If  $i = k$ , then  $x$  is the answer - it is larger than  $k - 1 = i - 1$  elements.

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We are looking for the third smallest element ( $i = 3$ )

And in our case the pivot is in the fourth position,  $k = 4$

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The third smallest element  
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For the same reason, if  $i > k$ , the answer cannot be in the first part.

# Running Time

Algorithm **Selection**( $\mathbf{A}[1, \dots, n]$ ,  $i$ )

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$T(|S_{\max}|)$

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Before we proceed, we have to bound  $|S_{\max}|$ .



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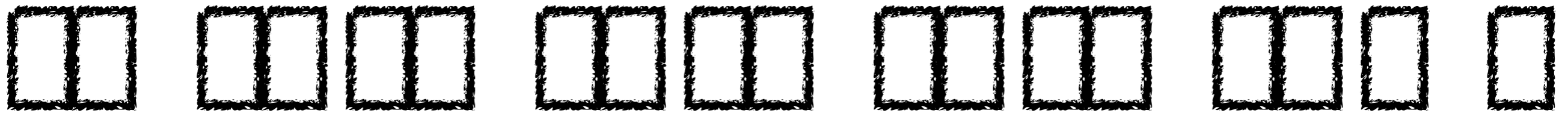
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Except possibly the group containing  $x$  and the group that has fewer than 5 elements.

# Median of medians



# Bounding the size of the subarrays

What is the total number of elements larger than  $x$ ?

$$3 \left( \left\lfloor \frac{1}{2} \cdot \left\lfloor \frac{n}{5} \right\rfloor \right\rfloor - 2 \right) \geq \frac{3n}{10} - 6$$

# elements  $> x$  in each of those groups (points to 3)  
 # groups with "baby medians"  $> x$  (points to  $\lfloor \frac{1}{2} \cdot \lfloor \frac{n}{5} \rfloor \rfloor$ )  
 # groups (points to  $\lfloor \frac{n}{5} \rfloor$ )  
 # groups who could be exceptions (points to  $-2$ )

This means that the size of the lower subarray is at most

$$7n/10 + 6$$

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Back to the recurrence:

$$T(n) \leq T(n/5) + T(|S_{\max}|) + bn = T(n/5) + T(7n/10 + 6) + bn$$

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$$\begin{aligned} T(n) &\leq c(n/5) + c(7n/10 + 6) + bn \\ &= 9cn/10 + 6c + bn \\ &= cn + (-cn/10 + 6c + bn) \end{aligned}$$

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This is at most  $cn$  whenever  $-cn/10 + 6c + bn \leq 0$ , or equivalently, when  $c \geq 10bn/(n - 60)$ .

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This is at most  $cn$  whenever  $-cn/10 + 6c + bn \leq 0$ , or equivalently, when  $c \geq 10bn/(n - 60)$ .

If  $n \geq 120$ , then  $n/(n - 60) \leq 2$  and then, it suffices to have  $c \geq 20b$ .



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Let  $a = \max\{ T(n) / n , n \leq 120\}$  and let  $c = \max\{a, 20b\}$ .

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We will prove the statement by induction.

# Solving the recurrence

We want to show that there is some constant  $c > 0$ , such that  $T(n) \leq cn$  for all  $n > 0$ .

Let  $a = \max\{ T(n) / n , n \leq 120\}$  and let  $c = \max\{a, 20b\}$ .

We will prove the statement by induction.

**Base case:** For every  $n \leq 120$ ,  $T(n) \leq \max\{ T(n) / n , n \leq 120\} n$   
 $= an \leq \max\{a, 20b\}n = cn$

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**Inductive Step:** Suppose that it holds for all  $n$  up to  $k = 120$ . Then for  $n = k + 1$ , we have  $T(n) \leq cn + (-cn/10 + 6c + bn)$

This follows from the previous slide and the fact that  $n > 120$  and  $c \geq 20b$ .