

# **Algorithms and Data Structures**

Matrix Chain Multiplication (cont), Longest Common  
Subsequence

# Dynamic Programming

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Solve the subproblems from the smallest to the largest. When you solve a subproblem, **store the solution** (e.g., in an array) and use it to solve the larger subproblems.

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RECTANGULAR-MATRIX-MULTIPLY( $A, B, C, p, q, r$ )
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1  for  $i = 1$  to  $p$ 
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```
2      for  $j = 1$  to  $r$ 
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3          for  $k = 1$  to  $q$ 
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The running time is the number of scalar multiplications, i.e.,  $pqr$ .

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If we do  $((A_1 \cdot A_2) \cdot A_3)$  we have  $10 \cdot 100 \cdot 5 = 5000$  scalar multiplications +  $10 \cdot 5 \cdot 50 = 2500$  scalar multiplications for a total of  $7500$ .



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If instead we do  $(A_1 \cdot (A_2 \cdot A_3))$  we have  $100 \cdot 5 \cdot 50 = 25000$  scalar multiplications +  $10 \cdot 100 \cdot 50 = 50000$  scalar multiplications for a total of  $75000$ .

# Full Parenthesisation

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We can think of the input as a sequence of dimensions  $\langle p_0, p_1, p_2, \dots, p_n \rangle$ .

# A Dynamic Programming algorithm (bottom-up)

MATRIX-CHAIN-ORDER( $p, n$ )

```
1  let  $m[1:n, 1:n]$  and  $s[1:n - 1, 2:n]$  be new tables
2  for  $i = 1$  to  $n$  // chain length 1
3       $m[i, i] = 0$ 
4  for  $l = 2$  to  $n$  //  $l$  is the chain length
5      for  $i = 1$  to  $n - l + 1$  // chain begins at  $A_i$ 
6           $j = i + l - 1$  // chain ends at  $A_j$ 
7           $m[i, j] = \infty$ 
8          for  $k = i$  to  $j - 1$  // try  $A_{i:k}A_{k+1:j}$ 
9               $q = m[i, k] + m[k + 1, j] + p_{i-1}p_kp_j$ 
10             if  $q < m[i, j]$ 
11                  $m[i, j] = q$  // remember this cost
12                  $s[i, j] = k$  // remember this index
13  return  $m$  and  $s$ 
```

# Computing a solution

We have the following relation:

$$M[i, j] = \begin{cases} 0, & \text{if } i = j, \\ \min_{k \in [i, j)} \{M[i, k] + M[k + 1, j] + p_{i-1}p_k p_j\} & \text{if } i < j. \end{cases}$$

We could now define a recursive algorithm straightforwardly using this.

```
RECURSIVE-MATRIX-CHAIN(p, i, j)
1  if i == j
2      return 0
3  m[i, j] = ∞
4  for k = i to j - 1
5      q = RECURSIVE-MATRIX-CHAIN(p, i, k)
           + RECURSIVE-MATRIX-CHAIN(p, k + 1, j)
           + pi-1pkpj
6      if q < m[i, j]
7          m[i, j] = q
8  return m[i, j]
```

# Memoization

We write the procedure recursively in a natural manner, but we modify it to save the result of each subproblem in a table.

The procedure first checks if we have computed the solution already in our table.

If we have, it uses it without calling itself recursively.

If we have not, then it recurses to compute it.

# A Dynamic Programming algorithm (top-down)

MEMOIZED-MATRIX-CHAIN( $p, n$ )

```
1 let  $m[1:n, 1:n]$  be a new table
2 for  $i = 1$  to  $n$ 
3   for  $j = i$  to  $n$ 
4      $m[i, j] = \infty$ 
5 return LOOKUP-CHAIN( $m, p, 1, n$ )
```

LOOKUP-CHAIN( $m, p, i, j$ )

```
1 if  $m[i, j] < \infty$ 
2   return  $m[i, j]$ 
3 if  $i == j$ 
4    $m[i, j] = 0$ 
5 else for  $k = i$  to  $j - 1$ 
6    $q =$  LOOKUP-CHAIN( $m, p, i, k$ )
       + LOOKUP-CHAIN( $m, p, k + 1, j$ ) +  $p_{i-1}p_kp_j$ 
7   if  $q < m[i, j]$ 
8      $m[i, j] = q$ 
9 return  $m[i, j]$ 
```



# Longest Common Subsequence

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**Motivation:** A strand of DNA consists of a string of bases, with possible values **A**denine, **C**ytosine, **G**uanine, and **T**hymine.

We can express a strand of DNA as a string, e.g.,

$S_1 = \text{ACCGGTCGAGTGCGCGGAAGCCGGCCAA}$

$S_2 = \text{GTCGTTCGGAATGCCGTTGCTCTGTAAA}$

We would like to compare two strands to see how similar they are (in order to see how similar two organisms are).

Various way to do that. Here: Find the longest possible strand  $S_3$  which is a subsequence of both  $S_1$  and  $S_2$ .

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$Z = \langle z_1, z_2, \dots, z_k \rangle$  is a *subsequence* of  $X = \langle x_1, x_2, \dots, x_m \rangle$  if there exists a strictly increasing sequence  $\langle i_1, i_2, \dots, i_k \rangle$  of indices of  $X$  such that for all  $j = 1, 2, \dots, k$  we have  $x_{i_j} = z_j$ .

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Example:  $Z = \langle B, C, D, B \rangle$  is a subsequence of  $X = \langle A, B, C, B, D, A, B \rangle$ .

# Brute Force

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Exponential time.

# Optimal Substructure

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Given a sequence  $X = \langle x_1, x_2, \dots, x_m \rangle$ , the  $i$ th *prefix* of  $X$  (for  $i = 1, 2, \dots, m$ ) is defined as  $X_i = \langle x_1, x_2, \dots, x_i \rangle$ .

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**Example:**  $X = \langle A, B, C, B, D, A, B \rangle$ ,  $X_4 = \langle A, B, C, B \rangle$  and  $X_0 = \langle \rangle$ .

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Example:  $X = \langle A, B, C, B, D, A, B \rangle$ ,  $X_4 = \langle A, B, C, B \rangle$  and  $X_0 = \langle \rangle$ .

## Theorem (Optimal Substructure):

Let  $X = \langle x_1, x_2, \dots, x_m \rangle$  and  $Y = \langle y_1, y_2, \dots, y_n \rangle$  be two sequences, and let  $Z = \langle z_1, z_2, \dots, z_k \rangle$  be any **LCS** of  $X$  and  $Y$ . Then.

1. If  $x_m = y_n$ , then  $z_k = x_m = y_n$  and  $Z_{k-1}$  is an **LCS** of  $X_{m-1}$  and  $Y_{n-1}$ .
2. If  $x_m \neq y_n$  and  $z_k \neq x_m$ , then  $Z$  is an **LCS** of  $X_{m-1}$  and  $Y$ .
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Assume by contradiction that there exists a longer common subsequence  $W$  of  $X_{m-1}$  and  $Y_{n-1}$  that has length at least  $k$ .

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If we append  $x_m = y_n$  to  $W$  we obtain a common subsequence of  $X$  and  $Y$  of length  $k + 1$  contradicting the fact that  $Z$  is a LCS.

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Let  $X = \langle x_1, x_2, \dots, x_m \rangle$  and  $Y = \langle y_1, y_2, \dots, y_n \rangle$  be two sequences, and let  $Z = \langle z_1, z_2, \dots, z_k \rangle$  be any **LCS** of  $X$  and  $Y$ . Then.

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## Proof:

(2) Since  $z_k \neq x_m$ ,  $Z$  is a common subsequence of  $X_{m-1}$  and  $Y$ . Assume by contradiction that it is no a **LCS**.

# Optimal Substructure

## Theorem (Optimal Substructure):

Let  $X = \langle x_1, x_2, \dots, x_m \rangle$  and  $Y = \langle y_1, y_2, \dots, y_n \rangle$  be two sequences, and let  $Z = \langle z_1, z_2, \dots, z_k \rangle$  be any **LCS** of  $X$  and  $Y$ . Then.

1. If  $x_m = y_n$ , then  $z_k = x_m = y_n$  and  $Z_{k-1}$  is an **LCS** of  $X_{m-1}$  and  $Y_{n-1}$ .
2. If  $x_m \neq y_n$  and  $z_k \neq x_m$ , then  $Z$  is an **LCS** of  $X_{m-1}$  and  $Y$ .
3. If  $x_m \neq y_n$  and  $z_k \neq y_n$ , then  $Z$  is an **LCS** of  $X$  and  $Y_{n-1}$ .

## Proof:

(2) Since  $z_k \neq x_m$ ,  $Z$  is a common subsequence of  $X_{m-1}$  and  $Y$ . Assume by contradiction that it is no a **LCS**.

Let  $W$  be a longer common subsequence, of length at least  $k + 1$ .

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Let  $W$  be a longer common subsequence, of length at least  $k + 1$ .

But  $W$  is a common subsequence of  $X$  and  $Y$ , contradicting the fact that  $Z$  is a **LCS**.

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  - Find a **LCS** of  $X$  and  $Y_{n-1}$ .
  - Choose the one that is longer.

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$$C[i, j] = \begin{cases} 0, & \text{if } i = 0 \text{ or } j = 0, \\ C[i - 1, j - 1] + 1, & \text{if } i, j > 0 \text{ and } x_i = y_j, \\ \max\{C[i, j - 1], C[i - 1, j]\} & \text{if } i, j > 0 \text{ and } x_i \neq y_j. \end{cases}$$









# A Dynamic Programming Algorithm

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
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4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
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8 else PRINT-LCS( $b, X, i, j - 1$ )
```

0	0	0	0
0	★	★	★
0	★	★	
0			
0			

# A Dynamic Programming Algorithm

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$ 
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = "\searrow"$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = "\uparrow"$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = "\leftarrow"$ 
16 return  $c$  and  $b$ 
```

PRINT-LCS( $b, X, i, j$ )

```
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3 if  $b[i, j] == "\searrow"$ 
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0	0	0	0
0	★	★	★
0	★	★	★
0			
0			

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0	0	0	0
0	★	★	★
0	★	★	★
0			
0			

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```

0	0	0	0
0	★	★	★
0	★	★	★
0			
0			

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```

0	0	0	0
0	↑	★	★
0	★	★	★
0			
0			

# A Dynamic Programming Algorithm

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8 else PRINT-LCS( $b, X, i, j - 1$ )
```

0	0	0	0
0	★	★	★
0	★	★	★
0			
0			



# A Dynamic Programming Algorithm

LCS-LENGTH( $X, Y, m, n$ )

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```

0	0	0	0
0 ←	*	*	*
0	*	*	*
0			
0			

# A Dynamic Programming Algorithm

LCS-LENGTH( $X, Y, m, n$ )

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```

0	0	0	0
0	★	★	★
0	★	★	★
0			
0			

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0					
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0					
<i>B</i>	0					
<i>C</i>	0					
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# Example

LCS-LENGTH( $X, Y, m, n$ )

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0					
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$$C[0,1] = 0$$

# Example

LCS-LENGTH( $X, Y, m, n$ )

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```


	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0					
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$$C[0,1] = 0$$

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# Example

LCS-LENGTH( $X, Y, m, n$ )

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10       $b[i, j] = \text{"↖"}$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$  
13       $b[i, j] = \text{"↑"}$ 
14    else  $c[i, j] = c[i, j - 1]$ 
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0					
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
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```

	$B$	$D$	$C$	$A$	$B$	$A$
	0	0	0	0	0	0
$A$	0	0				
$B$	0					
$C$	0					
$B$	0					
$D$	0					
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$B$	0					

$$C[0,1] = 0$$

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# Example

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0				
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<i>C</i>	0					
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# Example

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13       $b[i, j] = \text{“}\uparrow\text{”}$ 
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```

	$B$	$D$	$C$	$A$	$B$	$A$
	0	0	0	0	0	0
$A$	0	0				
$B$	0					
$C$	0					
$B$	0					
$D$	0					
$A$	0					
$B$	0					

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# Example

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10       $b[i, j] = \nwarrow$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = \uparrow$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = \leftarrow$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0				
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$$C[0,1] = 0$$

$$C[1,0] = 0$$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```

1  let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2  for  $i = 1$  to  $m$ 
3       $c[i, 0] = 0$ 
4  for  $j = 0$  to  $n$ 
5       $c[0, j] = 0$ 
6  for  $i = 1$  to  $m$       // compute table entries in row-major order
7      for  $j = 1$  to  $n$ 
8          if  $x_i == y_j$       Is  $x_1 = y_2$ ?
9               $c[i, j] = c[i - 1, j - 1] + 1$ 
10              $b[i, j] = \text{"↖"}$ 
11         elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12              $c[i, j] = c[i - 1, j]$ 
13              $b[i, j] = \text{"↑"}$ 
14         else  $c[i, j] = c[i, j - 1]$ 
15              $b[i, j] = \text{"←"}$ 
16  return  $c$  and  $b$ 

```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$$C[0,1] = 0 \quad C[0,2] = 0$$

$$C[1,0] = 0$$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
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7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$  // is  $x_1 = y_2$ ?
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
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16 return  $c$  and  $b$ 
```

	$B$	$D$	$C$	$A$	$B$	$A$
	0	0	0	0	0	0
$A$	0	0				
$B$	0					
$C$	0					
$B$	0					
$D$	0					
$A$	0					
$B$	0					

$$C[0,1] = 0 \quad C[0,2] = 0$$

$$C[1,0] = 0 \quad C[1,1] = 0$$

# Example

LCS-LENGTH( $X, Y, m, n$ )

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8          if  $x_i == y_j$       Is  $x_1 = y_2$ ?
9               $c[i, j] = c[i - 1, j - 1] + 1$ 
10              $b[i, j] = "\nwarrow"$ 
11          elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12              $c[i, j] = c[i - 1, j]$  ←
13              $b[i, j] = "\uparrow"$ 
14          else  $c[i, j] = c[i, j - 1]$ 
15              $b[i, j] = "\leftarrow"$ 
16  return  $c$  and  $b$ 

```


	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$C[0,1] = 0$     $C[0,2] = 0$   
 $C[1,0] = 0$     $C[1,1] = 0$

# Example

LCS-LENGTH( $X, Y, m, n$ )

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13              $b[i, j] = "\uparrow"$ 
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0			
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$C[0,1] = 0$     $C[0,2] = 0$   
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

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# Example

LCS-LENGTH( $X, Y, m, n$ )

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8     if  $x_i == y_j$ 
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11     elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
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13        $b[i, j] = \text{“}\uparrow\text{”}$ 
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0			
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
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# Example

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0	0		
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
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$$C[0,1] = 0 \quad C[0,2] = 0$$

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# Example

LCS-LENGTH( $X, Y, m, n$ )

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8          if  $x_i == y_j$       Is  $x_1 = y_4$ ?
9               $c[i, j] = c[i - 1, j - 1] + 1$ 
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$C[0,1] = 0$     $C[0,2] = 0$   
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# Example

LCS-LENGTH( $X, Y, m, n$ )

```

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
<i>B</i>	0					

$$C[0,1] = 0 \quad C[0,2] = 0$$

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# Example

LCS-LENGTH( $X, Y, m, n$ )

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0	1		
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
<i>A</i>	0					
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# Example

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0	1		
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
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$C[0,1] = 0$     $C[0,2] = 0$   
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# Example

LCS-LENGTH( $X, Y, m, n$ )

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0	0	1	
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>	0	0	0	0	0	0
<i>B</i>	0					
<i>C</i>	0					
<i>B</i>	0					
<i>D</i>	0					
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# Example

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```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0						
<i>C</i>	0						
<i>B</i>	0						
<i>D</i>	0						
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12        $c[i, j] = c[i - 1, j]$ 
13        $b[i, j] = \text{“}\uparrow\text{”}$ 
14     else  $c[i, j] = c[i, j - 1]$ 
15          $b[i, j] = \text{“}\leftarrow\text{”}$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1				
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$ 
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = \text{“}\searrow\text{”}$ 
11     elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12        $c[i, j] = c[i - 1, j]$ 
13        $b[i, j] = \text{“}\uparrow\text{”}$ 
14     else  $c[i, j] = c[i, j - 1]$ 
15          $b[i, j] = \text{“}\leftarrow\text{”}$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1				
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						

Next to compute:  $C[3,3]$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$  // is  $x_3 = y_3$ ?
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = \text{“}\searrow\text{”}$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = \text{“}\uparrow\text{”}$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = \text{“}\leftarrow\text{”}$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1				
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						

Next to compute:  $C[3,3]$

# Example

```
LCS-LENGTH( $X, Y, m, n$ )
1  let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2  for  $i = 1$  to  $m$ 
3       $c[i, 0] = 0$ 
4  for  $j = 0$  to  $n$ 
5       $c[0, j] = 0$ 
6  for  $i = 1$  to  $m$            // compute table entries in row-major order
7      for  $j = 1$  to  $n$ 
8          if  $x_i == y_j$            Is  $x_3 = y_3$ ?
9               $c[i, j] = c[i - 1, j - 1] + 1$ 
10              $b[i, j] = \nwarrow$ 
11         elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12              $c[i, j] = c[i - 1, j]$ 
13              $b[i, j] = \uparrow$ 
14         else  $c[i, j] = c[i, j - 1]$ 
15              $b[i, j] = \leftarrow$ 
16  return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1				
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						

Next to compute:  $C[3,3]$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$  Is  $x_3 = y_3$ ?
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = \text{“}\searrow\text{”}$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = \text{“}\uparrow\text{”}$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = \text{“}\leftarrow\text{”}$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1				
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						



Next to compute:  $C[3,3]$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$  Is  $x_3 = y_3$ ?
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = \nwarrow$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = \uparrow$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = \leftarrow$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
	0	0	0	0	0	0	
<i>A</i>	0	0	0	0	1	1	1
<i>B</i>	0	1	1	1	1	2	2
<i>C</i>	0	1	1	2			
<i>B</i>	0						
<i>D</i>	0						
<i>A</i>	0						
<i>B</i>	0						

Next to compute:  $C[3,3]$

# Example

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$ 
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = \text{“}\swarrow\text{”}$ 
11     elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12        $c[i, j] = c[i - 1, j]$ 
13        $b[i, j] = \text{“}\uparrow\text{”}$ 
14     else  $c[i, j] = c[i, j - 1]$ 
15          $b[i, j] = \text{“}\leftarrow\text{”}$ 
16 return  $c$  and  $b$ 
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
	0	0	0	0	0	0
<i>A</i>	0	0	0	0	1	1
<i>B</i>	0	1	1	1	1	2
<i>C</i>	0	1	1	2	2	2
<i>B</i>	0	1	1	2	2	3
<i>D</i>	0	1	2	2	2	3
<i>A</i>	0	1	2	2	3	3
<i>B</i>	0	1	2	2	3	4



# Example

```
PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```
PRINT-LCS( $b, X, i, j$ )
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6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```
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8  else PRINT-LCS( $b, X, i, j - 1$ )
```

		<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```

PRINT-LCS( $b, X, i, j$ )
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5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```

		<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```
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8  else PRINT-LCS( $b, X, i, j - 1$ )
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```

PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
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3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```

		<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```
PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
```

	<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>	
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑

# Example

```

PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```

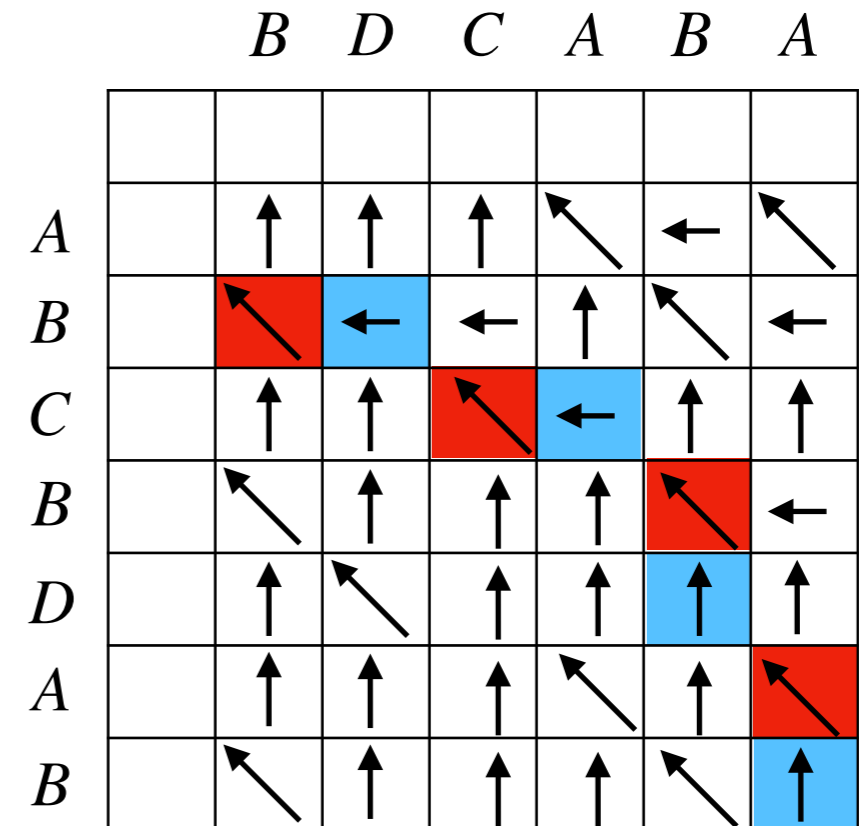
		<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>A</i>		↑	↑	↑	↖	←	↖
<i>B</i>		↖	←	←	↑	↖	←
<i>C</i>		↑	↑	↖	←	↑	↑
<i>B</i>		↖	↑	↑	↑	↖	←
<i>D</i>		↑	↖	↑	↑	↑	↑
<i>A</i>		↑	↑	↑	↖	↑	↖
<i>B</i>		↖	↑	↑	↑	↖	↑



# Example

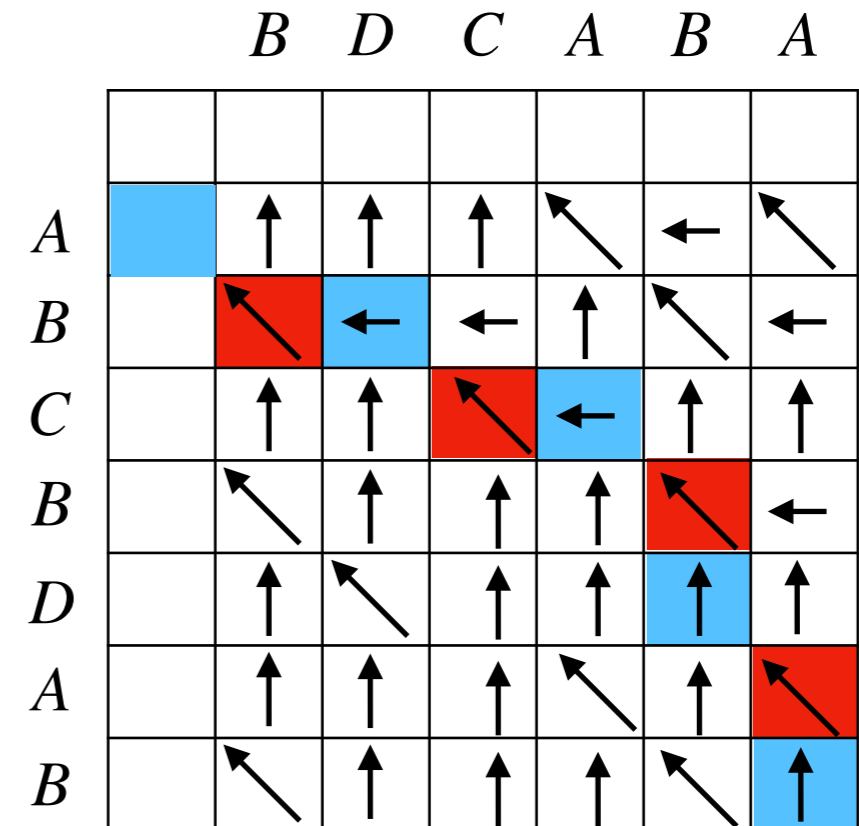
```

PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```



# Example

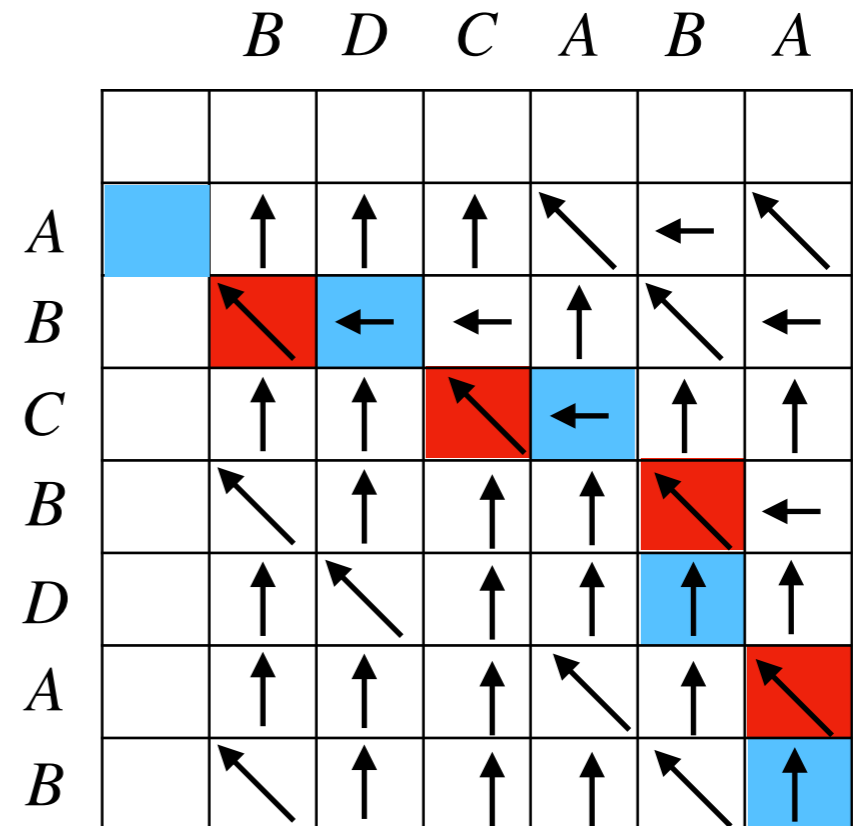
```
PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
```



# Example

```

PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```

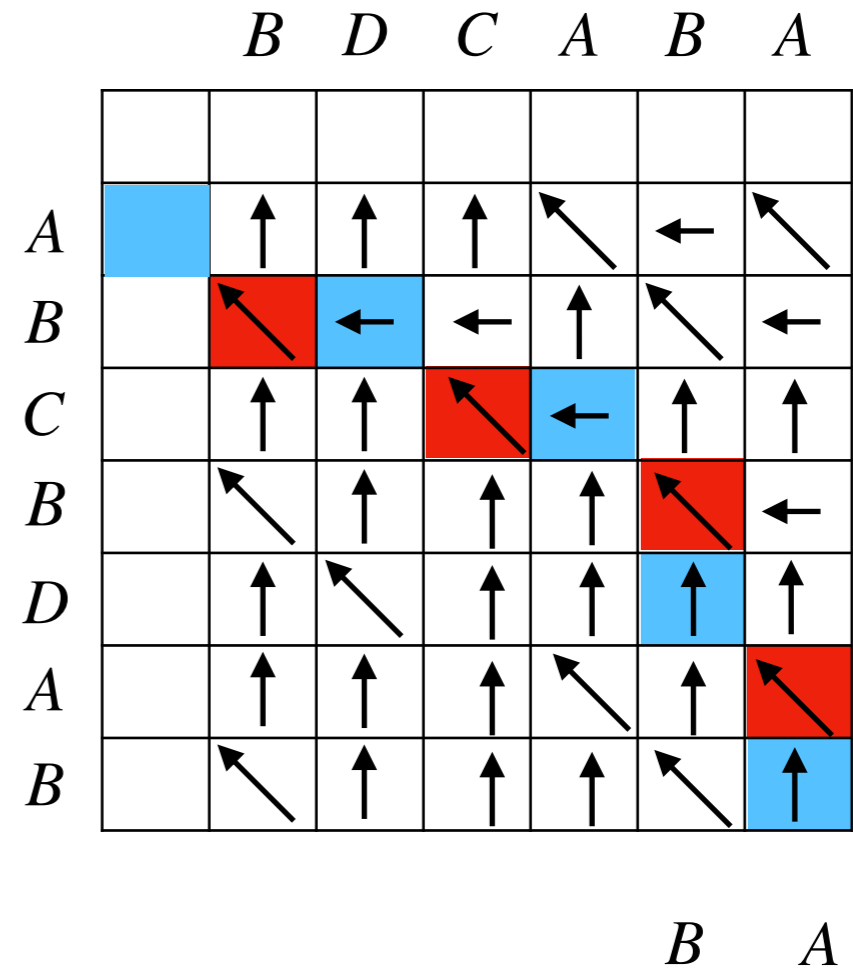


*A*

# Example

```

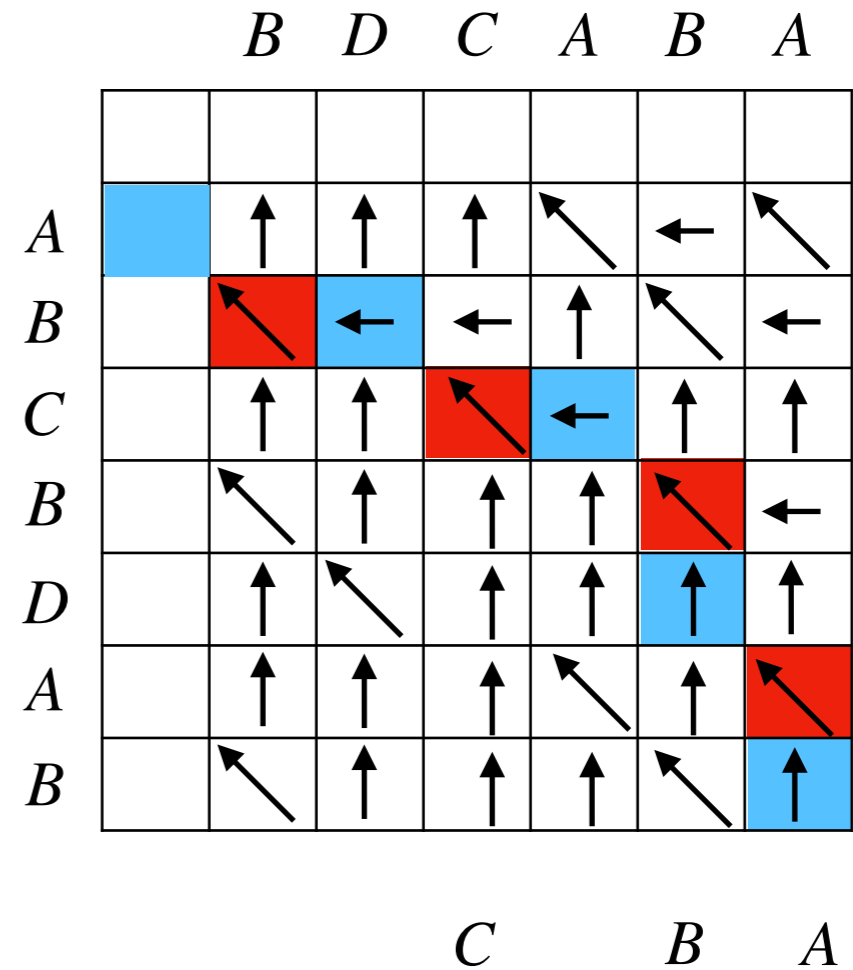
PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
    
```



# Example

```

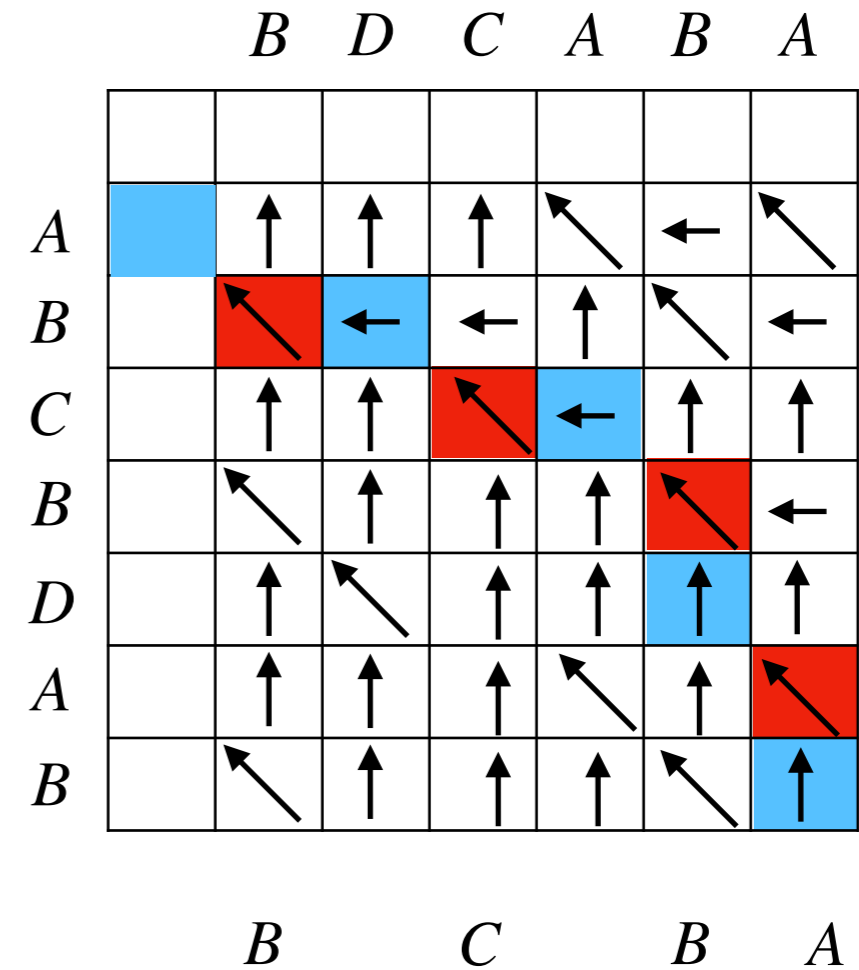
PRINT-LCS( $b, X, i, j$ )
1  if  $i == 0$  or  $j == 0$ 
2      return // the LCS has length 0
3  if  $b[i, j] == "\searrow"$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$  // same as  $y_j$ 
6  elseif  $b[i, j] == "\uparrow"$ 
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```



# A Dynamic Programming Algorithm

LCS-LENGTH( $X, Y, m, n$ )

```
1 let  $b[1:m, 1:n]$  and  $c[0:m, 0:n]$  be new tables
2 for  $i = 1$  to  $m$ 
3    $c[i, 0] = 0$ 
4 for  $j = 0$  to  $n$ 
5    $c[0, j] = 0$ 
6 for  $i = 1$  to  $m$  // compute table entries in row-major order
7   for  $j = 1$  to  $n$ 
8     if  $x_i == y_j$ 
9        $c[i, j] = c[i - 1, j - 1] + 1$ 
10       $b[i, j] = "\nwarrow"$ 
11    elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
12       $c[i, j] = c[i - 1, j]$ 
13       $b[i, j] = "\uparrow"$ 
14    else  $c[i, j] = c[i, j - 1]$ 
15       $b[i, j] = "\leftarrow"$ 
16 return  $c$  and  $b$ 
```

PRINT-LCS( $b, X, i, j$ )

```
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2   return // the LCS has length 0
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7   PRINT-LCS( $b, X, i - 1, j$ )
8 else PRINT-LCS( $b, X, i, j - 1$ )
```

0	0	0	0
0	*	*	*
0	*	*	*
0			
0			

# A Dynamic Programming Algorithm

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```

PRINT-LCS( $b, X, i, j$ )

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```

} do we need this?

0	0	0	0
0	*	*	*
0	*	*	*
0			
0			



# A Dynamic Programming Algorithm

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5       $c[0, j] = 0$ 
6  for  $i = 1$  to  $m$            // compute table entries in row-major order
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14          else  $c[i, j] = c[i, j - 1]$ 
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```

can we use less space here?

0	0	0	0
0	*	*	*
0	*	*	*
0			
0			

PRINT-LCS( $b, X, i, j$ )

```

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