Advanced Database Systems (ADBS), University of Edinburgh, 2023/24

Conjunctive Queries: Evaluation and Static Analysis

(Chapter 14 and 15 of DBT)

[DBT] Database Theory, https://github.com/pdm-book/community

Semantics of Conjunctive Queries

- A match of a conjunctive query $Q(x_1,\ldots,x_k)$:- body in a database D is a homomorphism h from the set of atoms body to the set of atoms D
- The answer to $Q(x_1,\ldots,x_k)$:- body over D is the set of k-tuples

  $Q(D) := \{(h(x_1),\ldots,h(x_k)) \mid h \text{ is a match of } Q \text{ in } D\}$
- The answer consists of the witnesses for the distinguished variables of Q

Pattern Matching Problem

List the airlines that fly directly from London to Glasgow

{
Flight(VIE,LHR,BA),
Flight(LHR,EDI,BA),
Flight(LGW,GLA,U2),
Flight(LCA,VIE,OS),
Airport(VIE,Vienna),
Airport(LHR,London),
Airport(LGW,London),
Airport(LCA,Larnaca),
Airport(GLA,Glasgow),
Airport(EDI,Edinburgh)
}

Q(z) :- Airport(x,London), Airport(y,Glasgow), Flight(x,y,z)

Pattern Matching Problem

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Airport(GLA,Glasgow),
Airport(EDI,Edinburgh)
}

$\{x \mapsto LGW,\ y \mapsto GLA,\ z \mapsto U2,$
$London \mapsto London,\ Glasgow \mapsto Glasgow\}$

Q(z) :- Airport(x,London), Airport(y,Glasgow), Flight(x,y,z)

Query Evaluation

- Understand the complexity of evaluating a conjunctive query over a database
- What to measure? Queries may have a large output, and it would be misleading to count the output as “complexity”
- We therefore consider the following decision problem for **CQ**

> **CQ**-Evaluation
> **Input:** a database D, a CQ $Q(x_1,\dots,x_k)$:- body, and a tuple $(a_1,\dots,a_k)$ of values
> **Question:** $(a_1,\dots,a_k) \in Q(D)$?

combined complexity

Data Complexity of Query Evaluation

- Measures the complexity in terms of the size of the database - the query is fixed
- Meaningful in practice since the database is usually much bigger than the query
- We consider the following decision problem for a fixed CQ $Q(x_1,\dots,x_k)$:- body

> Q-Evaluation
> **Input:** a database D, and a tuple $(a_1,\dots,a_k)$ of values
> **Question:** $(a_1,\dots,a_k) \in Q(D)$?

Complexity of Query Evaluation

Theorem: CQ-Evaluation is NP-complete, and in PTIME in data complexity

Proof:

(NP-membership) Guess-and-check:

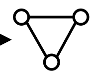
- Consider a database D, a CQ $Q(x_1,\dots,x_k)$:- body, and a tuple $(a_1,\dots,a_k)$ of values
- Guess a substitution h : terms(body) → terms(D)
- Verify that h is a match of Q in D, i.e., h(body) ⊆ D and $(h(x_1),\dots,h(x_k)) = (a_1,\dots,a_k)$

(NP-hardness) Reduction from 3-colorability

NP-hardness

(NP-hardness) Reduction from 3-colorability

> 3COL
> **Input:** an undirected graph **G** = (V,E)
> **Question:** is there a function c : V → {R,G,B} such that (v,u) ∈ E ⇒ c(v) ≠ c(u)?

Lemma: G is 3-colorable iff **G** can be mapped to $\mathbf{K_3}$, i.e., $\mathbf{G} \xrightarrow{\text{hom}}$ 

therefore, **G** is 3-colorable iff there is a match of Q_G in D = {E(x,y),E(y,z),E(z,x)}

the Boolean CQ that represents **G**

Complexity of Query Evaluation

Theorem: CQ-Evaluation is NP-complete, and in PTIME in data complexity

Proof:

(NP-membership) Guess-and-check:

- Consider a database D, a CQ $Q(x_1,\dots,x_k)$:- body, and a tuple $(a_1,\dots,a_k)$ of values
- Guess a substitution h : terms(body) → terms(D)
- Verify that h is a match of Q in D, i.e., h(body) ⊆ D and $(h(x_1),\dots,h(x_k)) = (a_1,\dots,a_k)$

(NP-hardness) Reduction from 3-colorability

(in PTIME) For every substitution h : terms(body) → terms(D), check if h(body) ⊆ D and $(h(x_1),\dots,h(x_k)) = (a_1,\dots,a_k)$

Static Analysis

CQ-Satisfiability
Input: a conjunctive query Q
Question: is there a database D such that Q(D) is non-empty?

- If the answer is no, then the input query Q makes no sense
- **CQ**-Evaluation becomes trivial - the answer is always NO!

Static Analysis

CQ-Equivalence
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \equiv Q_2$? or $Q_1(D) = Q_2(D)$ for every database D?

- Replace a query Q_1 with a query Q_2 that is easier to evaluate
- But, we have to be sure that $Q_1(D) = Q_2(D)$ for every database D

Static Analysis

CQ-Containment
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \subseteq Q_2$? or $Q_1(D) \subseteq Q_2(D)$ for every database D?

- Equivalence boils down to two containment checks
- Clearly, $Q_1(D) = Q_2(D)$ iff $Q_1(D) \subseteq Q_2(D)$ and $Q_2(D) \subseteq Q_1(D)$

Complexity of Static Analysis

CQ-Satisfiability
Input: a conjunctive query Q
Question: is there a database D such that Q(D) is non-empty?

CQ-Equivalence
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \equiv Q_2$? or $Q_1(D) = Q_2(D)$ for every database D?

CQ-Containment
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \subseteq Q_2$? or $Q_1(D) \subseteq Q_2(D)$ for every database D?

Canonical Database

- Convert a conjunctive query Q into a database D[Q] - the canonical database of Q
- Given a conjunctive query of the form Q(**x**) :- body, D[Q] is obtained from body by replacing each variable x with a new value c(x) = $\underline{x}$
- E.g., given Q(x,y) :- R(x,y), P(y,z,w), R(z,x), then D[Q] = {R($\underline{x}$,$\underline{y}$), P($\underline{y}$,$\underline{z}$,$\underline{w}$), R($\underline{z}$,$\underline{x}$)}
- **Note:** The mapping c : {variables in body} → {new values} is a bijection, where c(body) = D[Q] and c^{-1}(D[Q]) = body

Satisfiability of CQs

CQ-Satisfiability
Input: a conjunctive query Q
Question: is there a database D such that Q(D) is non-empty?

Theorem: A conjunctive query Q is always satisfiable

Proof: Due to its canonical database - Q(D[Q]) is trivially non-empty

Equivalence and Containment of CQs

CQ-Equivalence
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \equiv Q_2$? or $Q_1(D) = Q_2(D)$ for every database D?

CQ-Containment
Input: two conjunctive queries Q_1 and Q_2
Question: $Q_1 \subseteq Q_2$? or $Q_1(D) \subseteq Q_2(D)$ for every database D?

$Q_1 \equiv Q_2$ iff $Q_1 \subseteq Q_2$ and $Q_2 \subseteq Q_1$

$Q_1 \subseteq Q_2$ iff $Q_1 \equiv (Q_1 \wedge Q_2)$

...thus, we can safely focus on **CQ**-Containment

Homomorphism Theorem

A query homomorphism from $Q_1(x_1,\dots,x_k)$:- $body_1$ to $Q_2(y_1,\dots,y_k)$:- $body_2$ is a substitution $h : terms(body_1) \rightarrow terms(body_2)$ such that:

1. h is a homomorphism from $body_1$ to $body_2$
2. $(h(x_1),\dots,h(x_k)) = (y_1,\dots,y_k)$

Homomorphism Theorem: Let Q_1 and Q_2 be conjunctive queries. It holds that:

$Q_1 \subseteq Q_2$ iff there exists a query homomorphism from Q_2 to Q_1

Homomorphism Theorem: Example

$$\begin{tikzcd}[column sep=small] Q_1(x,y) \text{ :- } & R(x,z), & S(z,z), & R(z,y) \\ Q_2(a,b) \text{ :- } \arrow[u] & R(a,c), \arrow[u] & S(c,d), \arrow[u] & R(d,b) \arrow[u] \end{tikzcd}$$

$h = \{a \mapsto x, b \mapsto y, c \mapsto z, d \mapsto z\}$

- h is a query homomorphism from Q_2 to Q_1 $\Rightarrow$ $Q_1 \subseteq Q_2$
- But, there is no homomorphism from Q_1 to Q_2 $\Rightarrow$ $Q_1 \subset Q_2$

Homomorphism Theorem: Proof

Assume that $Q_1(x_1,\dots,x_k)$:- $body_1$ and $Q_2(y_1,\dots,y_k)$:- $body_2$

($\Rightarrow$) $Q_1 \subseteq Q_2$ $\Rightarrow$ there exists a query homomorphism from Q_2 to Q_1

- Clearly, $(c(x_1),\dots,c(x_k)) \in Q_1(D[Q_1])$ - recall that $D[Q_1] = c(body_1)$
- Since $Q_1 \subseteq Q_2$, we conclude that $(c(x_1),\dots,c(x_k)) \in Q_2(D[Q_1])$
- Therefore, there exists a homomorphism h such that $h(body_2) \subseteq D[Q_1] = c(body_1)$ and $h((y_1,\dots,y_k)) = (c(x_1),\dots,c(x_k))$
- By construction, $c^{-1}(c(body_1)) = body_1$ and $c^{-1}((c(x_1),\dots,c(x_k))) = (x_1,\dots,x_k)$
- Therefore, $c^{-1} \circ h$ is a query homomorphism from Q_2 to Q_1

$$\begin{tikzcd} Q_2(y_1,\dots,y_k) \text{ :- } \arrow[d] & body_2 \arrow[d, "h"] \arrow[dd, bend left=60, "c^{-1} \circ h"] \\ Q_1(c(x_1),\dots,c(x_k)) \text{ :- } \arrow[d] & c(body_1) \arrow[d, "c^{-1}"] \\ Q_1(x_1,\dots,x_k) \text{ :- } & body_1 \end{tikzcd}$$

Homomorphism Theorem: Proof

Assume that $Q_1(x_1,\dots,x_k)$:- $body_1$ and $Q_2(y_1,\dots,y_k)$:- $body_2$

($\Leftarrow$) $Q_1 \subseteq Q_2$ $\Leftarrow$ there exists a query homomorphism from Q_2 to Q_1

- Consider a database D, and a tuple **t** such that $\mathbf{t} \in Q_1(D)$
- We need to show that $\mathbf{t} \in Q_2(D)$
- Clearly, there exists a homomorphism g such that $g(body_1) \subseteq D$ and $g((x_1,\dots,x_k)) = \mathbf{t}$
- By hypothesis, there exists a query homomorphism h from Q_2 to Q_1
- Therefore, $g(h(body_2)) \subseteq D$ and $g(h((y_1,\dots,y_k))) = \mathbf{t}$, which implies that $\mathbf{t} \in Q_2(D)$

$$\begin{tikzcd} Q_2(y_1,\dots,y_k) \text{ :- } \arrow[d] & body_2 \arrow[d, "h"] \arrow[dd, bend left=60, "g \circ h"] \\ Q_1(x_1,\dots,x_k) \text{ :- } \arrow[d] & body_1 \arrow[d, "g"] \\ \mathbf{t} & D \end{tikzcd}$$

Existence of a Query Homomorphism

Theorem: Let Q_1 and Q_2 be conjunctive queries. The problem of deciding whether there exists a query homomorphism from Q_2 to Q_1 is NP-complete

Proof:

(NP-membership) Guess a substitution, and verify that is a query homomorphism

(NP-hardness) Easy reduction from **CQ**-Evaluation

By applying the homomorphism theorem we get that:

Corollary: CQ-Equivalence and **CQ**-Containment are NP-complete