# Elements of Programming Languages Tutorial 7: Small-step semantics and type soundness

Week 9 (November 11–15, 2024)

Exercises marked  $\star$  are more advanced. Please try all unstarred exercises before the tutorial meeting.

# 1. Imperative programming

Write evaluation derivations for the following imperative programs, starting with the environment  $\sigma = [x = 3, y = 4]$ .

- (a) y := x + x
- (b) if x == y then x := x + 1 else y := y + 2
- (c) ( $\star$ ) while x < y do x := x + 1

## 2. Comparing large-step and small-step derivations

Write both large-step and small-step derivations for the following expressions. For the small-step derivations, construct the derivations of each  $e \mapsto e'$  step explicitly.

- (a)  $(\lambda x.x + 1) 42$
- (b)  $(\lambda x.if x == 1 then 2 else x + 1) 42$

### 3. Small-step derivations that go wrong

For each of the following expressions, show the small-step evaluation leading to the point where evaluation becomes stuck due to a dynamic type error. (There is no need to show the derivations of each step.)

- (a)  $((\lambda x.\lambda y. \text{let } z = x + y \text{ in } z + 1) 42)$  true
- (b)  $(\lambda x. \text{if } x \text{ then } x+1 \text{ else } x+2) \text{ true}$

# 4. Small-step rules for L<sub>Data</sub>

Recall that we defined the semantics for L<sub>Data</sub> using big-step rules, as follows:

$$e \Downarrow v$$

$$\frac{e_1 \Downarrow v_1 \quad e_2 \Downarrow v_2}{(e_1,e_2) \Downarrow (v_1,v_2)} \quad \frac{e \Downarrow (v_1,v_2)}{\mathsf{fst} \, e \Downarrow v_1} \quad \frac{e \Downarrow (v_1,v_2)}{\mathsf{snd} \, e \Downarrow v_2}$$
 
$$\frac{e \Downarrow v}{\mathsf{left}(e) \Downarrow \mathsf{left}(v)} \quad \frac{e \Downarrow \mathsf{left}(v_1) \quad e_1[v_1/x] \Downarrow v}{\mathsf{case} \, e \, \mathsf{of} \, \{\mathsf{left}(x) \Rightarrow e_1 \, ; \, \mathsf{right}(y) \Rightarrow e_2\} \Downarrow v}$$
 
$$\frac{e \Downarrow v}{\mathsf{right}(e) \Downarrow \mathsf{right}(v)} \quad \frac{e \Downarrow \mathsf{right}(v_2) \quad e_2[v_2/y] \Downarrow v}{\mathsf{case} \, e \, \mathsf{of} \, \{\mathsf{left}(x) \Rightarrow e_1 \, ; \, \mathsf{right}(y) \Rightarrow e_2\} \Downarrow v}$$

- (a) For each construct, write out equivalent small-step rules. Are there any design choices in translating the big-step rules to small-step rules?
- (b) (\*) Construct small-step derivations reducing the following expressions to values:
  - i.  $(\lambda p.(\operatorname{snd} p,\operatorname{fst} p+2))$  (17,42)
  - ii.  $(\lambda x.\mathtt{case}\ x\ \mathtt{of}\ \{\mathtt{left}(y).y+1\ ;\ \mathtt{right}(z).\ z\})\ (\mathtt{left}(42))$

# 5. $(\star)$ Type soundness for nondeterminism

This question builds on the *nondeterministic choice* construct mentioned in an earlier tutorial, with the following typing rules:

$$\Gamma \vdash e : \tau$$

$$\frac{\Gamma \vdash e_1 : \tau \quad \Gamma \vdash e_2 : \tau}{\Gamma \vdash e_1 \Box e_2 : \tau}$$

and small-step evaluation rules:

$$e \mapsto e'$$

$$\overline{e_1 \square e_2 \mapsto e_1} \qquad \overline{e_1 \square e_2 \mapsto e_2}$$

- (a) State the *preservation* property. Outline how we could prove the cases of preservation for nondeterministic expressions.
- (b) State the *progress* property. Outline how we could prove the cases of progress for nondeterministic expressions.