Introduction to Databases

(INFR10080)

(Deductive Databases)

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(Fall 2025)

THE UNIVERSITY of EDINBURGH

Changelog

v25.0 Initial version

Datalog rules (1)

$$\underbrace{H(\bar{x})}_{\text{head}} \;:-\; \underbrace{S_1(\bar{x}_1), \ldots, S_n(\bar{x}_n)}_{\text{body}}$$

Head predicate over variables and constants

Body conjunction of possibly negated atoms (**subgoals**)

- relational atoms (predicates)
- comparisons between variables/constants

- Head variables are implicitly **universally quantified**
- Body variables not in head are **existentially quantified**
- Rule: Body $\rightarrow$ Head (if the body is true, the head is true)

Datalog rules (2)

Example

$$\text{Lucky}(x) \;:-\; \text{Customer}(x, y, z), \text{Account}(u, z, x, w), w > 10000$$

In relational calculus we could write:

$$\text{Lucky} = \{\, x \mid \exists y, z, u, w\, \text{Customer}(x, y, z) \\ \wedge \text{Account}(u, z, x, w) \wedge w > 10000 \,\}$$

Safety

Every variable (in head or body)
appears in at least one non-negated relational atom

Not safe: $\text{BigNumber}(x) \;:-\; x > 1000000000$

Datalog programs

Program = set of Datalog rules

Example

$$\text{Parent}(x, y) \;:-\; \text{Mother}(x, y)$$
$$\text{Parent}(x, y) \;:-\; \text{Father}(x, y)$$

edb (extensional database) relations **stored** in the database
- can appear only in the body of rules

idb (intensional database) derived relations
- can appear both in the head or body of rules

From relational algebra to Datalog

Every relational algebra expression can be translated into Datalog

Projection $\pi_{\#2}(R)$

$$E(x) \;:-\; R(y, x)$$

Selection $\sigma_{\#1\,\mathbf{op}\,c}(R)$

$$E(x, y) \;:-\; R(x, y), x\,\mathbf{op}\,c$$

Product $R \times S$

$$E(x, y, w, z) \;:-\; R(x, y), S(w, z)$$

Difference $R - S$

$$E(x, y) \;:-\; R(x, y), \neg S(x, y)$$

Union $R \cup S$

$$E(x, y) \;:-\; R(x, y)$$
$$E(x, y) \;:-\; S(x, y)$$

From relational algebra to Datalog

Let R and S be relations over A, B

$$E = \underbrace{\pi_{A,D}\Big(\underbrace{\sigma_{B=C}\big(\underbrace{R \times \rho_{A\to C,\, B\to D}(S)}_{E_1}\big)}_{E_2}\Big)}_{E_3} \cup \underbrace{\rho_{B\to D}(R - S)}_{E_4}$$

$$\begin{aligned}
E_1(x,u,w,y) &:- R(x,u), S(w,y) \\
E_2(x,u,w,y) &:- E_1(x,u,w,y), u = w \\
E_3(x,y) &:- E_2(x,u,w,y) \\
E_4(x,y) &:- R(x,y), \neg S(x,y) \\
E(x,y) &:- E_3(x,y) \\
E(x,y) &:- E_4(x,y)
\end{aligned}$$

Limitations of relational algebra/calculus

Parent = table of pairs x, y where x is the parent of y

$$\begin{aligned}
\text{Parent} &= \{x, y \mid \text{Parent}(x,y)\} \\
\text{Grandparent} &= \{x, y \mid \exists z\ \text{Parent}(x,z) \wedge \text{Parent}(z,y)\} \\
\text{Great-grandparent} &= \{x, y \mid \exists z\ \text{Grandparent}(x,z) \wedge \text{Parent}(z,y)\}
\end{aligned}$$

For a **given** k, we can express the query Ancestor^k

But

We **cannot express** the Ancestor relation itself
that is: an Ancestor^k query that works for **every** k

Recursion in Datalog

The head relation of a rule can appear in its body

$$\text{Ancestor}(x, y) \;:-\; \text{Parent}(x, y)$$
$$\text{Ancestor}(x, y) \;:-\; \text{Ancestor}(x, z), \text{Parent}(z, y)$$

Intuition

x is an ancestor of y if

x is a parent of y

or

x is an ancestor of a parent of y

Dependency graph

IDB predicate P **depends on** (IDB) predicate Q
if there is a rule with P in the head and Q in a subgoal

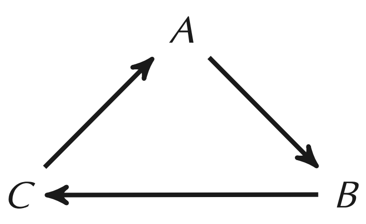
Dependency graph

nodes IDB predicates

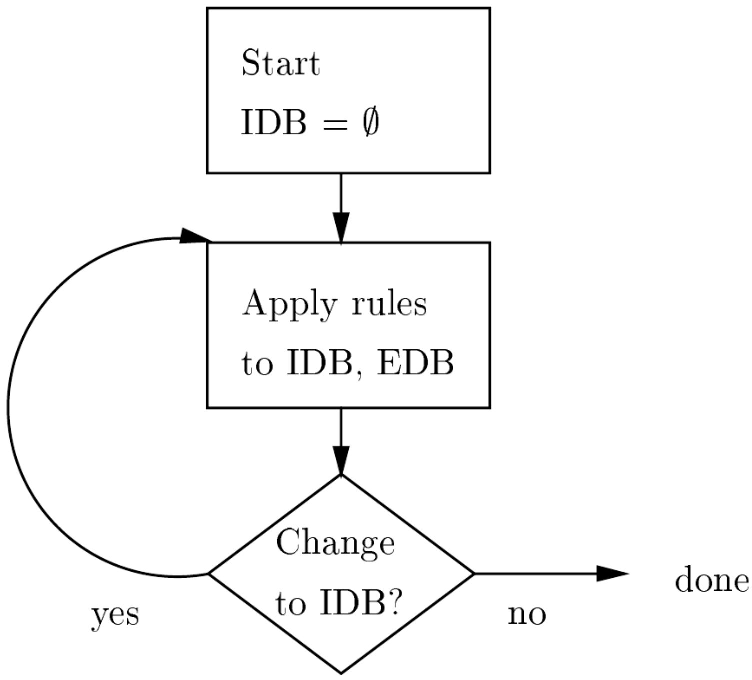
edges $P \to Q$ if P depends on Q

A **cycle** in the dependency graph means the program is recursive

$$C(x) \;:-\; A(y, x)$$
$$C(x) \;:-\; S(x, y), y > 1$$
$$B(x, y) \;:-\; C(x), P(x, y)$$
$$A(x, y) \;:-\; B(y, x)$$



Iterative Fixpoint Evaluation



Evaluation of recursive programs

- The **Parent** relation (EDB) never changes
- The **Ancestor** relation (IDB) is initially empty

  $\text{Ancestor}_0 = \varnothing$

- At step $i + 1$ compute:

  $\text{Ancestor}_{i+1}(x, y) \;:\!-\; \text{Parent}(x, y)$
  $\text{Ancestor}_{i+1}(x, y) \;:\!-\; \text{Ancestor}_i(x, z), \text{Parent}(z, y)$

- Stop when a **fixpoint** is reached

  $\text{Ancestor}_{i+1} = \text{Ancestor}_i$

Evaluation of recursive programs

IDB:

Ancestor	
John	Mary
John	Jane
Jane	Louis
Mary	Linda
Louis	Mark
John	Linda
John	Louis
Jane	Mark
John	Mark

EDB:

Parent	
John	Mary
John	Jane
Jane	Louis
Mary	Linda
Louis	Mark

$$\text{Ancestor}(x, y) :- \text{Parent}(x, y)$$
$$\text{Ancestor}(x, y) :- \text{Ancestor}(x, z), \text{Parent}(z, y)$$

Recursion in SQL

Suppose we have a table Parent with attributes name, child

```
WITH RECURSIVE Ancestor(name, descendant) AS (
     SELECT *
     FROM   Parent
     UNION
     SELECT A.name, P.child
     FROM   Ancestor A, Parent P
     WHERE  A.descendant = P.name
)
SELECT * FROM Ancestor ;
```

The definition mimics the structure of the Datalog program

Nonlinear recursion

The head relation can appear **more than once** in its body

$$\text{Ancestor}(x, y) :- \text{Parent}(x, y)$$
$$\text{Ancestor}(x, y) :- \text{Ancestor}(x, z), \text{Ancestor}(z, y)$$

Intuition

x is an ancestor of y if

x is a parent of y

or

x is an **ancestor of an ancestor** of y

Nonlinear recursion in SQL

Not supported

```
WITH RECURSIVE Ancestor(name, descendant) AS (
    SELECT *
    FROM   Parent
    UNION
    SELECT A.name, A2.desc
    FROM   Ancestor A1, Ancestor A2
    WHERE  A1.descendant = A2.name
)
SELECT * FROM Ancestor ;
```

ERROR: recursive reference to query ”ancestor” must not appear more than once

Recursive programs with negation

Consider the program $P = \{R(x) \ {:}{-}\ S(x), \neg R(x)\}$

IDB:

R

EDB:

S
1
2

Step 0 IDB relation R is empty

Step 1 $R = \{1, 2\}$

Step 2 $R = \varnothing$

Step 3 $R = \{1, 2\}$

Step 4 $R = \varnothing$

... **No fixpoint!**
Iteration never ends

Stratification

Partition a program P into a sequence of subprograms $P_1, \ldots, P_n$

- Each subprogram defines one or more IDB relations
- If a relation S is used positively in the definition of R then S must be defined **earlier or simultaneously** with R
- If a relation S is used negatively in the definition of R then S must be defined **strictly before** R

Stratum 0 EDB relations

Stratum i IDB relations that depend
positively on relations in any stratum $j \leq i$
negatively on relations in any stratum $j < i$

Stratification

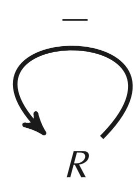
Stratum graph

nodes IDB predicates

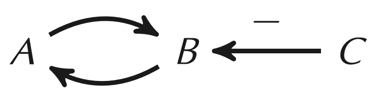
edges $P \rightarrow Q$ if P depends on Q
label the edge with "–" if Q is a negated subgoal

Stratifiable program: no cycle involving at least one negated edge

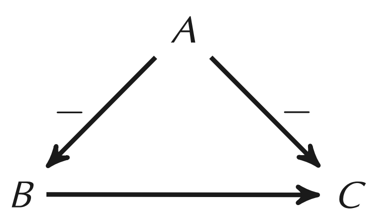
$R(x)$:– $S(x), \neg R(x)$



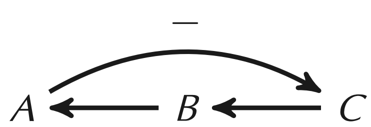
More examples



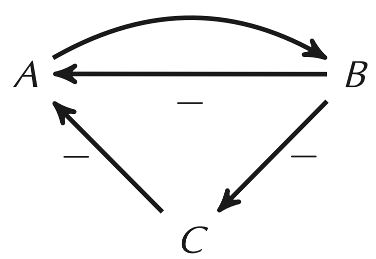
Stratifiable: $A = B < C$



Stratifiable: $C \leq B < A$



Not stratifiable: $A \leq B \leq C < A$



Not stratifiable: $A < B \leq A$

Stratified example

Which target nodes cannot be reached from any source node?

$$
\begin{aligned}
\mathrm{NoReach}(x) &:- \mathrm{Target}(x), \neg\mathrm{Reach}(x) && \text{(rule1)}\\
\mathrm{Reach}(x) &:- \mathrm{Source}(x) && \text{(rule2)}\\
\mathrm{Reach}(x) &:- \mathrm{Reach}(y), \mathrm{Link}(y, x) && \text{(rule3)}
\end{aligned}
$$

$$
\begin{tikzcd}
\mathrm{Reach} \arrow[loop above] & \mathrm{NoReach} \arrow[l, "-"']
\end{tikzcd}
$$

Stratum 0	Source, Link, Target
Stratum 1	Reach
Stratum 2	NoReach

Evaluation of stratified programs

P partitioned into a **sequence** $P_1, \ldots, P_n$

Gives us an order in which to apply (each group of) rules

At each iteration k, execute each subprogram in sequence

(1) Apply all the rules in P_1

(2) Apply all the rules in P_2

$\vdots$

(n) Apply all the rules in P_n

Evaluation of stratified programs

$$\begin{aligned}
\mathrm{Reach}(x) &:- \mathrm{Source}(x) && (P_1)\\
\mathrm{Reach}(x) &:- \mathrm{Reach}(y), \mathrm{Link}(y, x) && (P_1)\\
\mathrm{NoReach}(x) &:- \mathrm{Target}(x), \neg\mathrm{Reach}(x) && (P_2)
\end{aligned}$$

EDB

Source
1
2

Link	
1	2
3	4
2	4

Target
2
3
4

IDB

Reach
1
2
4

NoReach
3

Further remarks on SQL recursion

- Requires **stratified negation**
- Only **linear recursion**

Problems

Arithmetic operations
: introduce new values not present in the database

Multiset semantics
: rules must be applied several times
: **cycles in the data** must be detected