



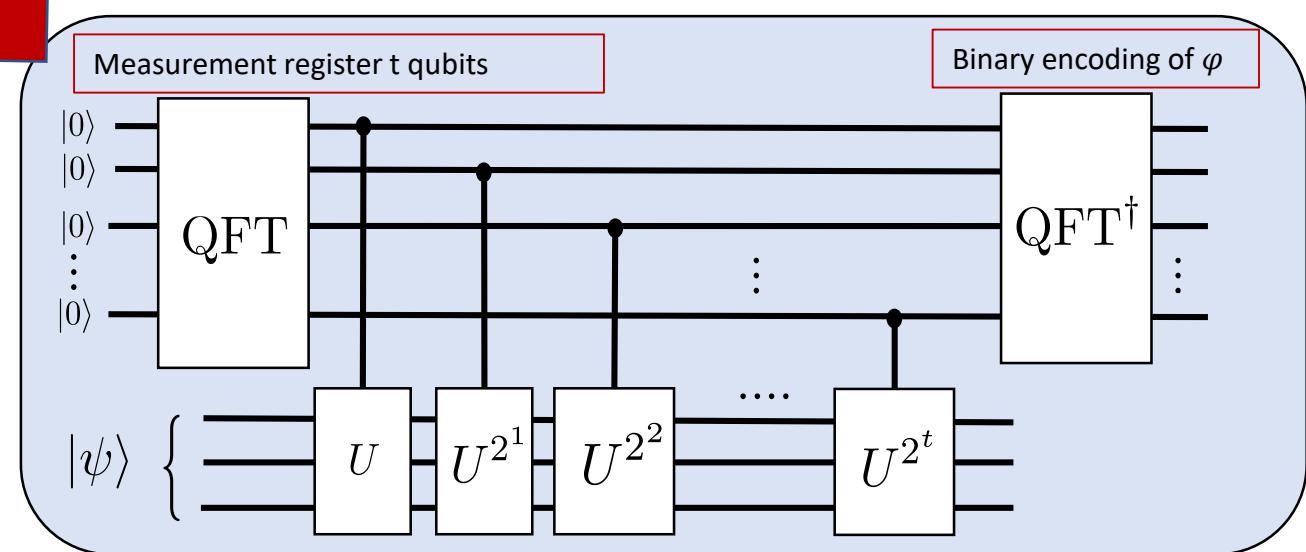
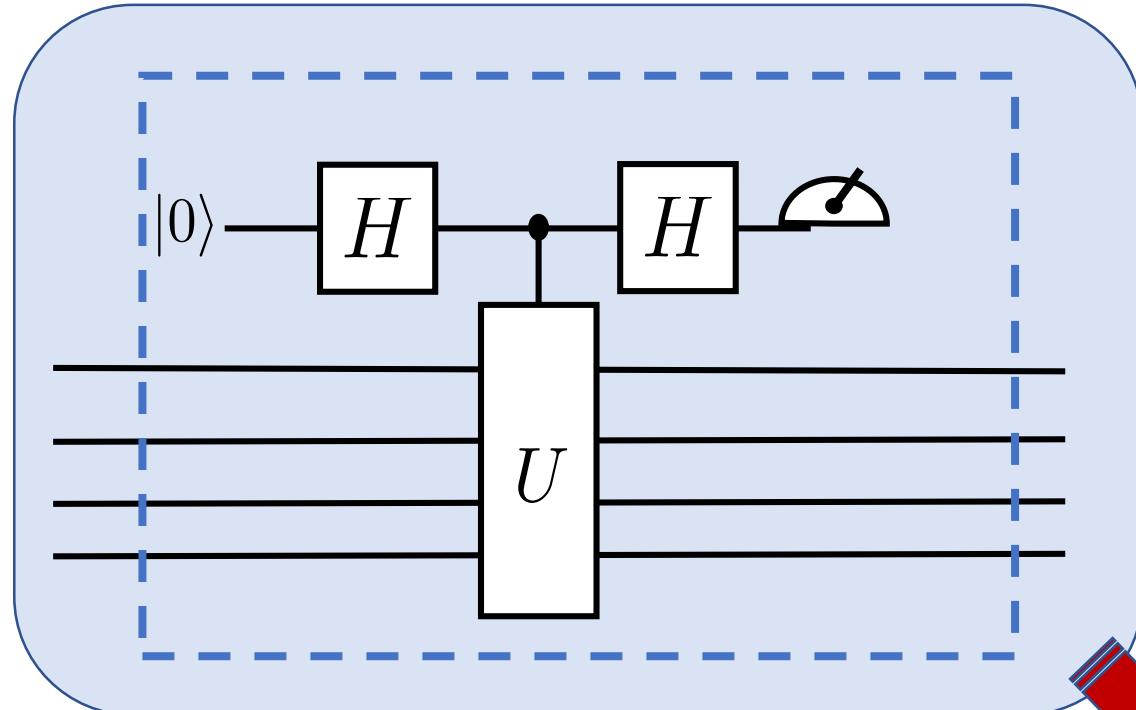
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Introduction to Quantum Computing

Lecture 18: Quantum Fourier Transform

Raul Garcia-Patron Sanchez

Generalization to N outcomes measurement





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Recap Hadamard Test



Hadamard Test – binary outcome measurements

We want to ask the question: "Are you in subspace Π_0 or Π_1 ?

$$\Pi_0 + \Pi_1 = I_{2^n}$$

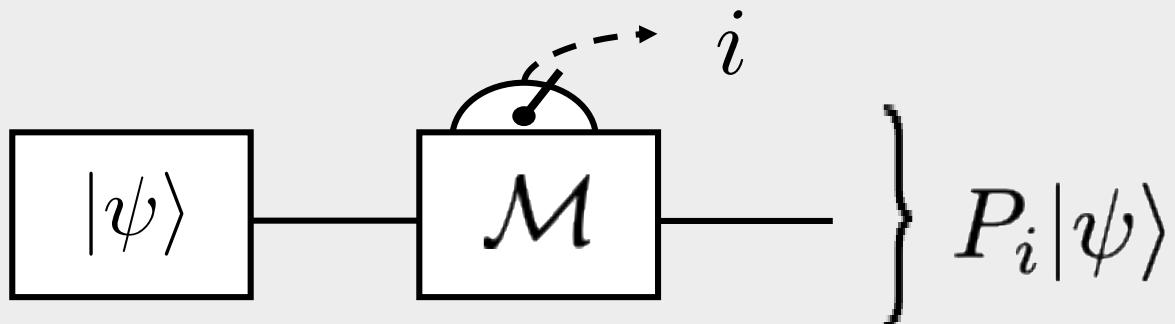
- Syndrome measurement in quantum error correction
- SWAP test: overlap between quantum states, measurement of purity with two copies

From the definition to designing a circuit implementing it

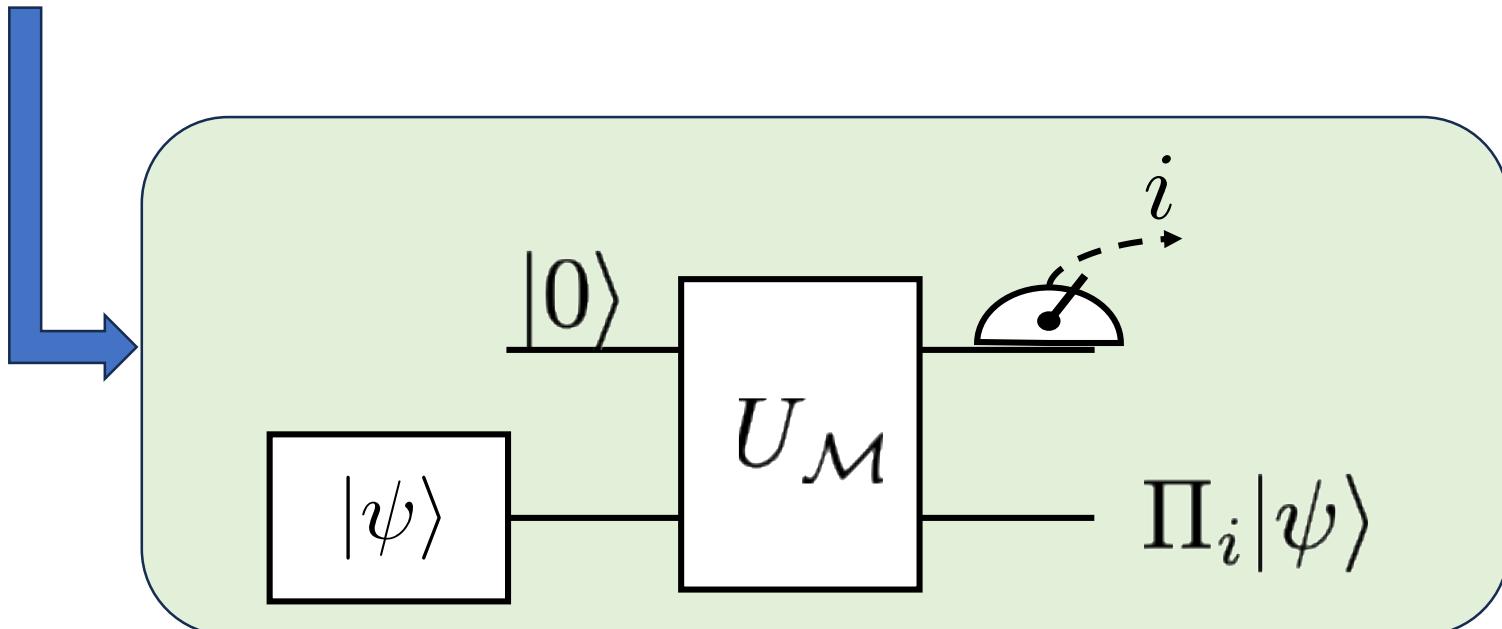
$$\Pi_0 + \Pi_1 = I_d$$

$$P(i) = ||\Pi_i|\psi\rangle||^2$$

Update: $\frac{\Pi_i|\psi\rangle}{||\Pi_i|\psi\rangle||}$

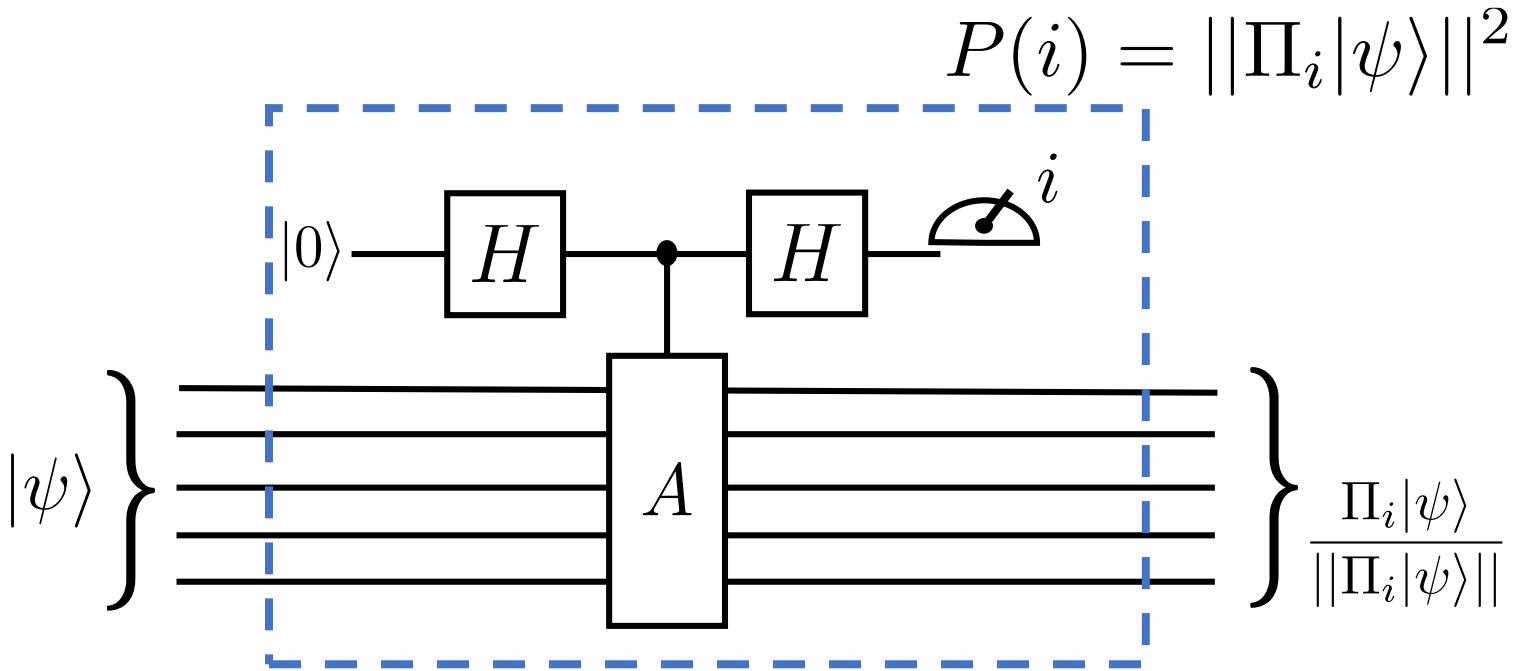


We need to find the circuit $U_{\mathcal{M}}$



General ± 1 eigenvalues measurement

$$\begin{aligned}\mathcal{H} &= \mathcal{H}_0 \oplus \mathcal{H}_1 \\ \Pi_0 + \Pi_1 &= I \\ A &= \Pi_0 - \Pi_1\end{aligned}$$



$$\begin{aligned}|0\rangle \otimes |\psi\rangle &\rightarrow |+\rangle \otimes |\psi\rangle \rightarrow |0\rangle \otimes |\psi\rangle + |1\rangle \otimes A|\psi\rangle \\ &= |0\rangle \otimes (\Pi_0 + \Pi_1)|\psi\rangle + |1\rangle \otimes (\Pi_0 - \Pi_1)|\psi\rangle \\ &= (|0\rangle + |1\rangle) \otimes \Pi_0|\psi\rangle + (|0\rangle - |1\rangle) \otimes \Pi_1|\psi\rangle \\ &\rightarrow |0\rangle \otimes \Pi_0|\psi\rangle + |1\rangle \otimes \Pi_1|\psi\rangle\end{aligned}$$

- Apply measurement on upper qubit: $\tilde{\Pi}_i = |i\rangle\langle i| \otimes I$

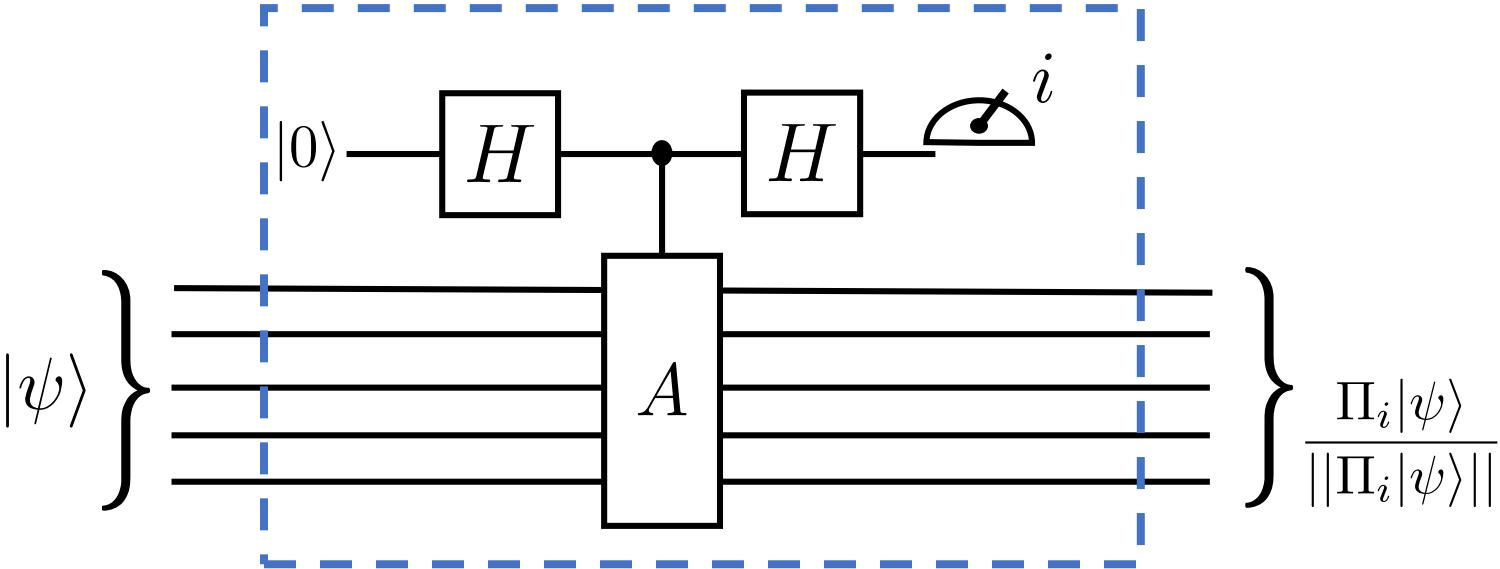
General ± 1 eigenvalues measurement

$$\mathcal{H} = \mathcal{H}_0 \oplus \mathcal{H}_1$$

$$\Pi_0 + \Pi_1 = I$$

$$A = \Pi_0 - \Pi_1$$

$$P(i) = \|\Pi_i|\psi\rangle\|^2$$



$$|0\rangle \otimes |\psi\rangle \rightarrow |0\rangle \otimes \Pi_0|\psi\rangle + |1\rangle \otimes \Pi_1|\psi\rangle$$

$$\tilde{\Pi}_0 = |0\rangle\langle 0| \otimes I$$

$$\tilde{\Pi}_0|\tilde{\phi}\rangle = (|0\rangle\langle 0| \otimes I)(|0\rangle \otimes \Pi_0|\psi\rangle + |1\rangle \otimes \Pi_1|\psi\rangle) = |0\rangle \otimes \Pi_0|\psi\rangle$$

$$P(0) = \|\tilde{\Pi}_0|\tilde{\phi}\rangle\|^2 = (\langle\tilde{\phi}|\tilde{\Pi}_0)(\tilde{\Pi}_0|\tilde{\phi}\rangle) = (\langle 0| \otimes \langle\psi|\Pi_0)(|0\rangle \otimes \Pi_0|\psi\rangle)$$

$$= \langle 0|0\rangle\langle\psi|\Pi_0^2|\psi\rangle = \langle 0|0\rangle\langle\psi|\Pi_0|\psi\rangle = \|\Pi_0|\psi\rangle\|^2$$

General ± 1 eigenvalues measurement

$$\mathcal{H} = \mathcal{H}_0 \oplus \mathcal{H}_1$$

$$\Pi_0 + \Pi_1 = I$$

$$A = \Pi_0 - \Pi_1$$

$$|\tilde{\phi}\rangle = |0\rangle \otimes \Pi_e |\psi\rangle + |1\rangle \otimes \Pi_o |\psi\rangle$$

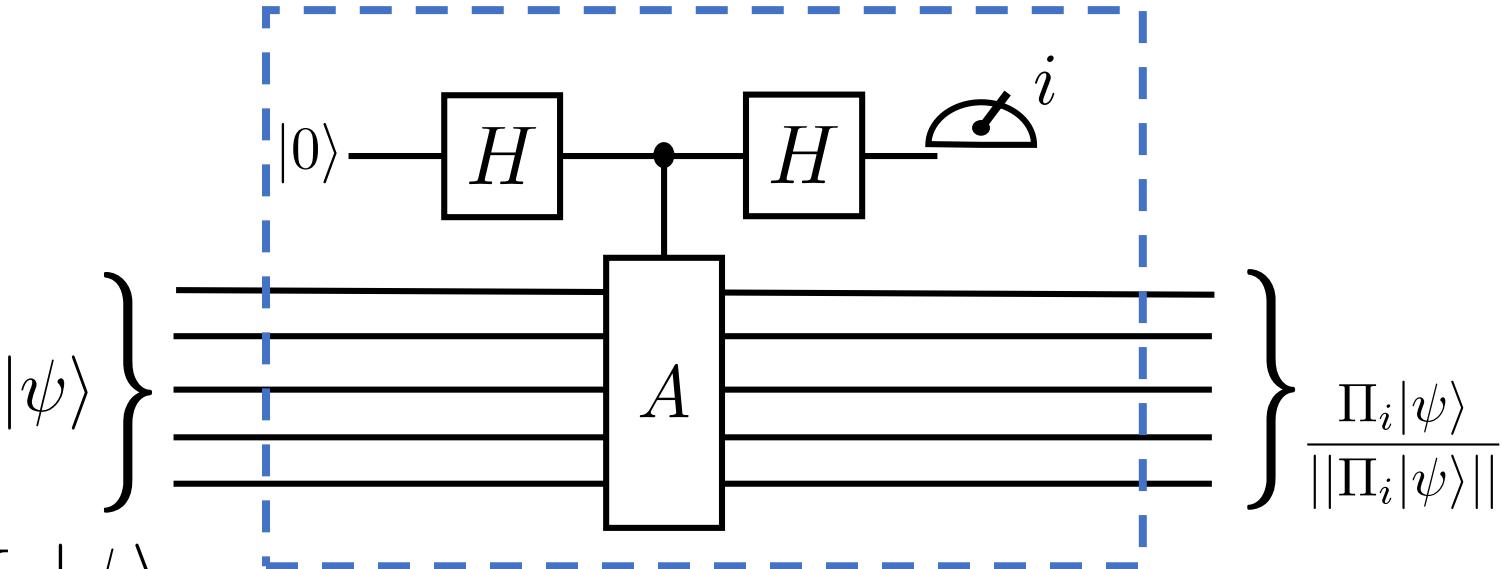
$$\tilde{\Pi}_0 = |0\rangle\langle 0| \otimes I$$

$$P(0) = \|\tilde{\Pi}_0|\tilde{\phi}\rangle\|^2 = \|\Pi_0|\psi\rangle\|^2$$

$$\tilde{\Pi}_1 = |1\rangle\langle 1| \otimes I$$

$$P(1) = \|\tilde{\Pi}_1|\tilde{\phi}\rangle\|^2 = \|\Pi_1|\psi\rangle\|^2$$

$$P(i) = \|\Pi_i|\psi\rangle\|^2$$



$$\frac{\tilde{\Pi}_0|\tilde{\phi}\rangle}{\|\tilde{\Pi}_0|\tilde{\phi}\rangle\|} = \frac{|0\rangle \otimes \Pi_0|\psi\rangle}{\|\Pi_0|\psi\rangle\|} = |0\rangle \otimes \frac{\Pi_0|\psi\rangle}{\|\Pi_0|\psi\rangle\|}$$

$$\frac{\tilde{\Pi}_1|\tilde{\phi}\rangle}{\|\tilde{\Pi}_1|\tilde{\phi}\rangle\|} = \frac{|1\rangle \otimes \Pi_1|\psi\rangle}{\|\Pi_1|\psi\rangle\|} = |1\rangle \otimes \frac{\Pi_1|\psi\rangle}{\|\Pi_1|\psi\rangle\|}$$



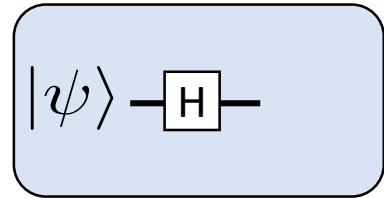
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Quantum Fourier Transform



From Hadamard to QFT

Hadamard Gate

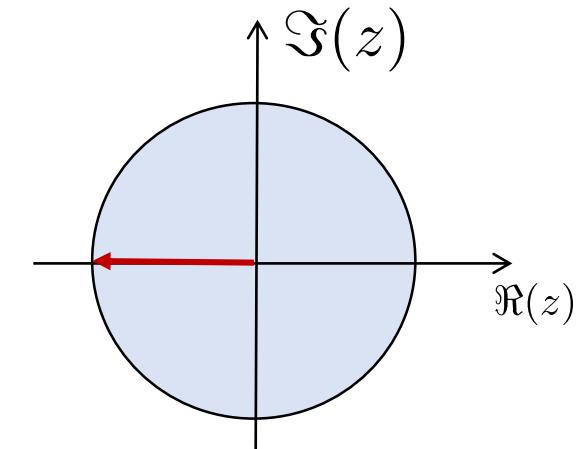


$$|x\rangle \xrightarrow{H} \frac{1}{\sqrt{2}} \sum_{y \in \{0,1\}} (-1)^{xy} |y\rangle$$

$$(-1)^x (-1)^y \equiv x \oplus y$$

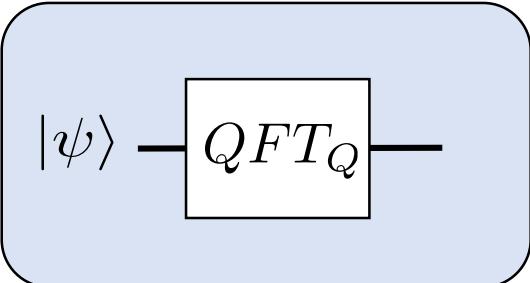
$$-1 = e^{i\pi}$$

$$(-1)^2 = 1$$



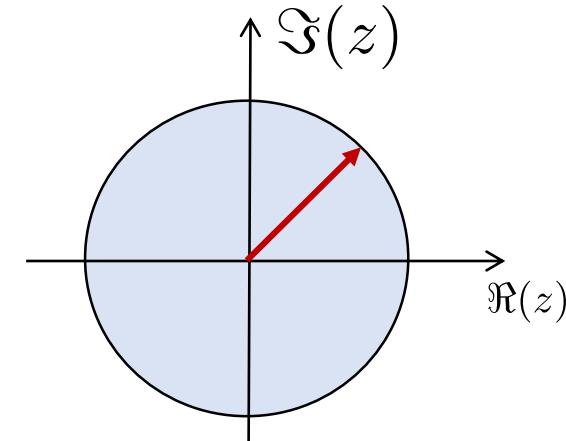
Fourier Transform over \mathbb{Z}_Q

$$\omega = e^{i2\pi/Q}$$



$$|x\rangle \xrightarrow{QFT_Q} \frac{1}{\sqrt{Q}} \sum_{y \in \mathbb{Z}_Q} \omega^{xy} |y\rangle$$

$$\omega^Q = 1$$



Basis of size N : $|0\rangle, |1\rangle, \dots, |Q-1\rangle$

$$\omega^x \omega^y = e^{ix2\pi/Q} e^{iy2\pi/Q} \equiv x + y \pmod{Q}$$

Quantum Fourier Transform over \mathbb{Z}_{2^n}

Binary representation of integers:

$$z \equiv z_1 z_2 \dots z_n$$

$$z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$$

Binary fraction:

$$0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$$

Quantum Fourier Transform over \mathbb{Z}_{2^n}

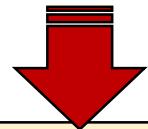
$$|x\rangle \xrightarrow{QFT_{2^n}} \frac{1}{2^{n/2}} \sum_{y=0}^{2^n-1} e^{2\pi i xy/2^n} |y\rangle$$

$$H|x_1\rangle = \frac{1}{\sqrt{2}}[|0\rangle + (-1)^{x_1}|1\rangle]$$

$$H|x_1\rangle = \frac{1}{2^{1/2}}[|0\rangle + e^{2\pi i 0.x_1}|1\rangle]$$

Binary representation: $z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$

Binary fraction: $0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$



We will now prove

$$x_1, x_2, \dots, x_{n-1}, x_n \rangle \rightarrow \frac{1}{2^{n/2}}(|0\rangle + e^{2\pi i 0.x_n}|1\rangle)(|0\rangle + e^{2\pi i 0.x_{n-1}x_n}|1\rangle) \dots (|0\rangle + e^{2\pi i 0.x_1x_2\dots x_n}|1\rangle)$$

Quantum Fourier Transform over \mathbb{Z}_{2^n}

$$|x\rangle \rightarrow \frac{1}{2^{n/2}} \sum_{y=0}^{2^n - 1} e^{2\pi i xy/2^n} |y\rangle$$

Binary representation: $z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$

Binary fraction: $0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$

$$= \frac{1}{2^{n/2}} \sum_{y_1=0}^1 \dots \sum_{y_n=0}^1 e^{2\pi i x (\sum_{l=1}^n y_l 2^{-l})} |y_1 \dots y_n\rangle$$

re-writing y in binary notation

$$= \frac{1}{2^{n/2}} \sum_{y_1=0}^1 \dots \sum_{y_n=0}^1 \bigotimes_{l=1}^n e^{2\pi i x y_l 2^{-l}} |y_l\rangle$$

exp. of sum = product of exp.

Property: $c(|u\rangle \otimes |v\rangle) = (c|u\rangle) \otimes |v\rangle = |u\rangle \otimes (c|v\rangle)$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[\sum_{y_l=0}^1 e^{2\pi i x y_l 2^{-l}} |y_l\rangle \right]$$

Property: $|u\rangle \otimes (|v_1\rangle + |v_2\rangle) = |u\rangle \otimes |v_1\rangle + |u\rangle \otimes |v_2\rangle$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i x 2^{-l}} |1\rangle \right]$$

Expanding each sum

Quantum Fourier Transform over \mathbb{Z}_{2^n}

$$|x\rangle \rightarrow \frac{1}{2^{n/2}} \sum_{y=0}^{2^n-1} e^{2\pi i xy/2^n} |y\rangle$$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[\sum_{y_l=0}^1 e^{2\pi i x y_l 2^{-l}} |y_l\rangle \right]$$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i x 2^{-l}} |1\rangle \right] = \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i 0.x_{n-l+1} \dots x_n} |1\rangle \right]$$

$$z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$$

$$e^{2\pi i z} = e^{2\pi i (z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n)} = 1 (z \in \mathbb{Z})$$

$$\frac{z}{2} = z_1 2^{n-2} + z_2 2^{n-3} + \dots + z_{n-2} 2 + z_{n-1} + \frac{z_n}{2} \Rightarrow e^{2\pi i \frac{z}{2}} = e^{2\pi i (z_1 2^{n-2} + z_2 2^{n-3} + \dots + z_{n-2} 2 + z_{n-1})} e^{2\pi i \frac{z_n}{2}} = e^{2\pi i 0.z_n}$$

$$\frac{z}{4} = z_1 2^{n-3} + z_2 2^{n-4} + \dots + z_{n-2} + \frac{z_{n-1}}{2} + \frac{z_n}{4} \Rightarrow e^{2\pi i \frac{z}{4}} = e^{2\pi i (z_1 2^{n-3} + z_2 2^{n-4} + \dots + z_{n-2})} e^{2\pi i (\frac{z_{n-1}}{2} + \frac{z_n}{4})} = e^{2\pi i 0.z_{n-1} z_n}$$

⋮

$$\frac{z}{2^l} = z_1 2^{n-1-l} + \dots + z_{n-l} + \frac{z_{n-l+1}}{2} + \dots + \frac{z_n}{2^l}$$

$$e^{2\pi i z 2^{-l}} = e^{2\pi i (z_1 2^{n-l-1} + z_2 2^{n-l-2} + \dots + z_{n-l})} e^{2\pi i 0.z_{n-l+1} z_{n-l+2} \dots z_n} = e^{2\pi i 0.z_{n-l+1} z_{n-l+2} \dots x_n}$$

Binary representation: $z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$

Binary fraction: $0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$

Quantum Fourier Transform over \mathbb{Z}_{2^n}

$$|x\rangle \rightarrow \frac{1}{2^{n/2}} \sum_{y=0}^{2^n-1} e^{2\pi i xy/2^n} |y\rangle$$

Binary representation: $z = z_1 2^{n-1} + z_2 2^{n-2} + \dots + z_{n-1} 2 + z_n$

Binary fraction: $0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$

$$= \frac{1}{2^{n/2}} \sum_{y_1=0}^1 \dots \sum_{y_n=0}^1 e^{2\pi i x (\sum_{l=1}^n y_l 2^{-l})} |y_1 \dots y_n\rangle$$

$$= \frac{1}{2^{n/2}} \sum_{y_1=0}^1 \dots \sum_{y_n=0}^1 \bigotimes_{l=1}^n e^{2\pi i x y_l 2^{-l}} |y_l\rangle$$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[\sum_{y_l=0}^1 e^{2\pi i x y_l 2^{-l}} |y_l\rangle \right]$$

$$= \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i x 2^{-l}} |1\rangle \right] = \frac{1}{2^{n/2}} \bigotimes_{l=1}^n \left[|0\rangle + e^{2\pi i 0.x_{n-l+1} \dots x_n} |1\rangle \right]$$

$$x_1, x_2, \dots, x_{n-1}, x_n \rangle \rightarrow \frac{1}{2^{n/2}} (|0\rangle + e^{2\pi i 0.x_n} |1\rangle) (|0\rangle + e^{2\pi i 0.x_{n-1} x_n} |1\rangle) \dots (|0\rangle + e^{2\pi i 0.x_1 x_2 \dots x_n} |1\rangle)$$

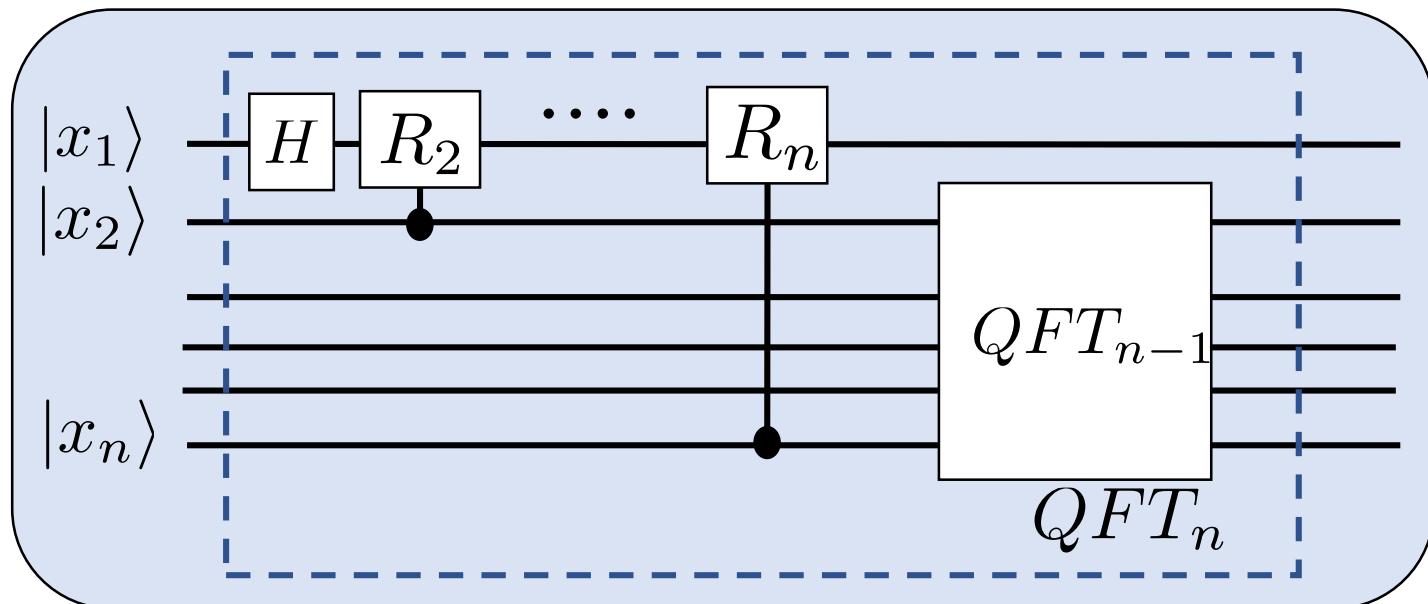
Quantum Fourier Transform over \mathbb{Z}_{2^n}

$$|x_1, x_2, \dots, x_{n-1}, x_n\rangle \rightarrow \frac{1}{2^{n/2}}(|0\rangle + e^{2\pi i 0.x_n}|1\rangle)(|0\rangle + e^{2\pi i 0.x_{n-1}x_n}|1\rangle) \dots (|0\rangle + e^{2\pi i 0.x_1x_2\dots x_n}|1\rangle)$$

$$R_k = \begin{bmatrix} 1 & 0 \\ 0 & e^{i2\pi/2^k} \end{bmatrix}$$

Binary fraction:

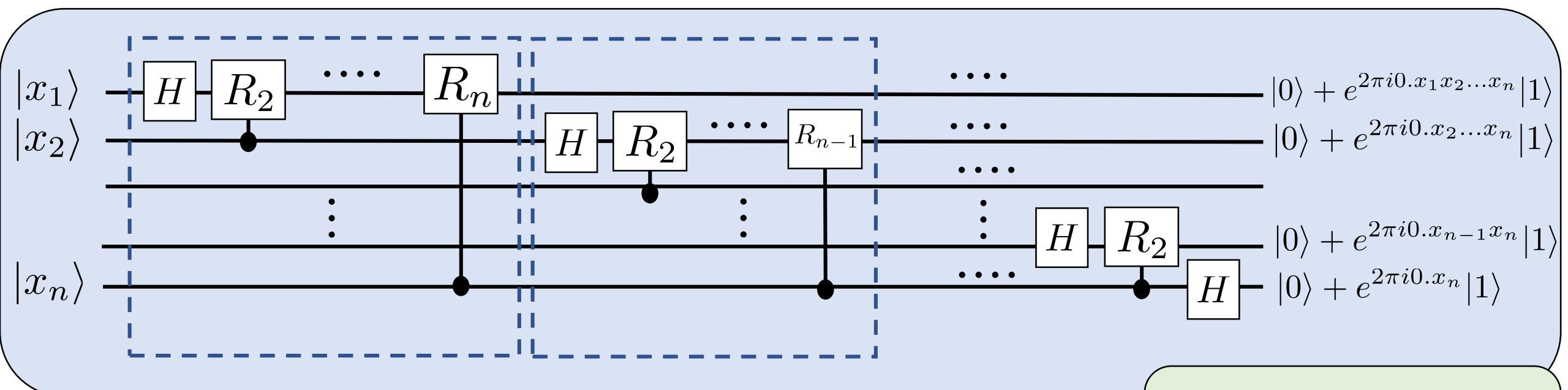
$$0.w_l w_{l+1} \dots w_m = \frac{w_l}{2} + \frac{w_{l+1}}{2^2} + \dots + \frac{w_m}{2^{m-l+1}}$$



$$\begin{aligned}
 |x_1\rangle \otimes |x_2 \dots x_n\rangle &\xrightarrow{H} \frac{1}{2^{1/2}} [|0\rangle + e^{2\pi i 0.x_1}|1\rangle] \otimes |x_2 \dots x_n\rangle \\
 &\xrightarrow{cR_2} \frac{1}{2^{1/2}} [|0\rangle + e^{2\pi i 0.x_1x_2}|1\rangle] \otimes |x_2 \dots x_n\rangle \\
 &\xrightarrow{cR_n} \frac{1}{2^{1/2}} [|0\rangle + e^{2\pi i 0.x_1x_2\dots x_n}|1\rangle] \otimes |x_2 \dots x_n\rangle
 \end{aligned}$$

Quantum Fourier Transform over \mathbb{Z}_{2^n}

$$|x_1, x_2, \dots, x_{n-1}, x_n\rangle \rightarrow \frac{1}{2^{n/2}}(|0\rangle + e^{2\pi i 0.x_n}|1\rangle) \dots (|0\rangle + e^{2\pi i 0.x_2\dots x_n}|1\rangle)(|0\rangle + e^{2\pi i 0.x_1 x_2 \dots x_n}|1\rangle)$$



- QFT up to a reverse of the order of qubits.

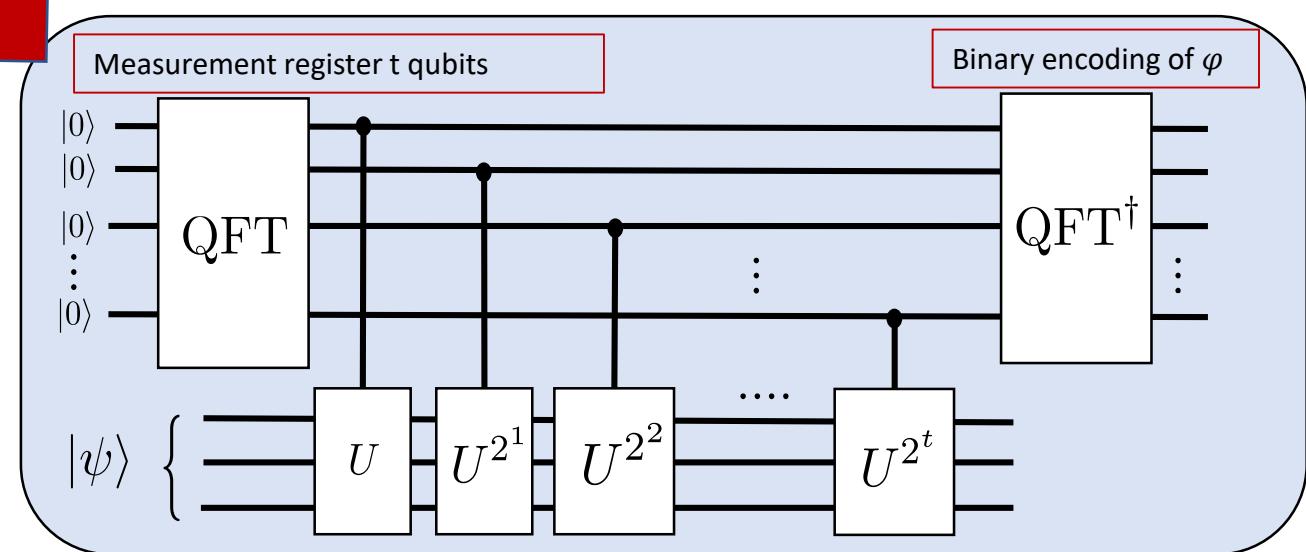
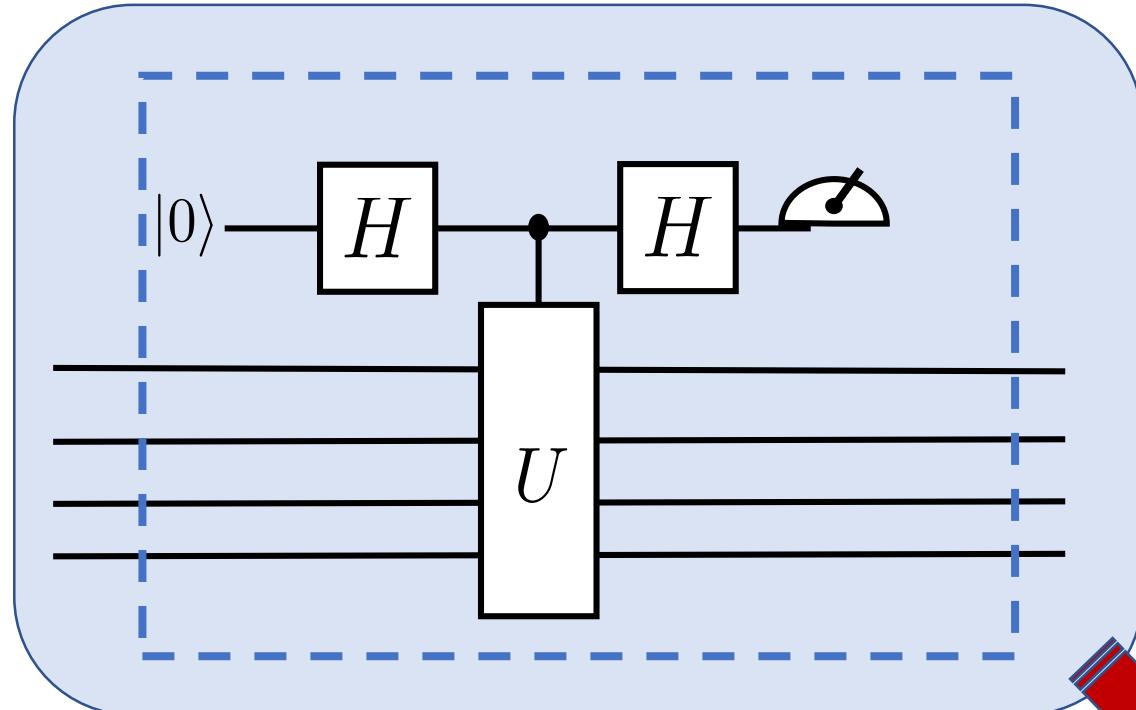
$$R_k = \begin{bmatrix} 1 & 0 \\ 0 & e^{i2\pi/2^k} \end{bmatrix}$$

- Unitary as composed of unitary gates.

- $n + (n - 1) + \dots + 1 = n(n + 1)/2$ gates are required
+ $n/2$ SWAP gates (3 CNOT)

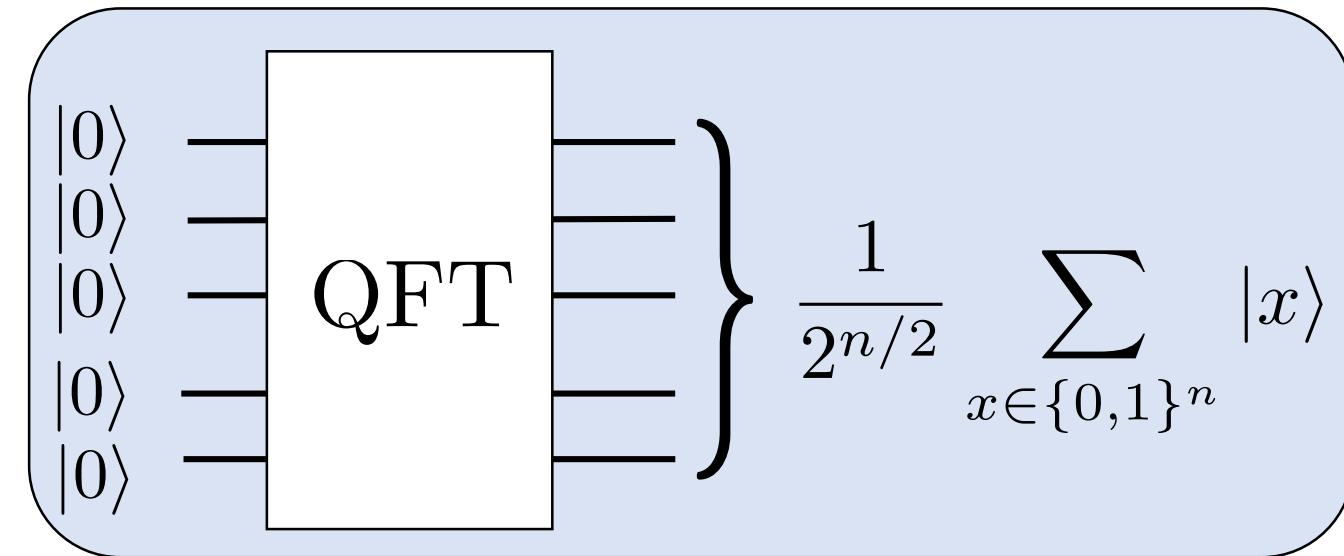
$\Theta(n^2)$
gates

Generalization to N outcomes measurement

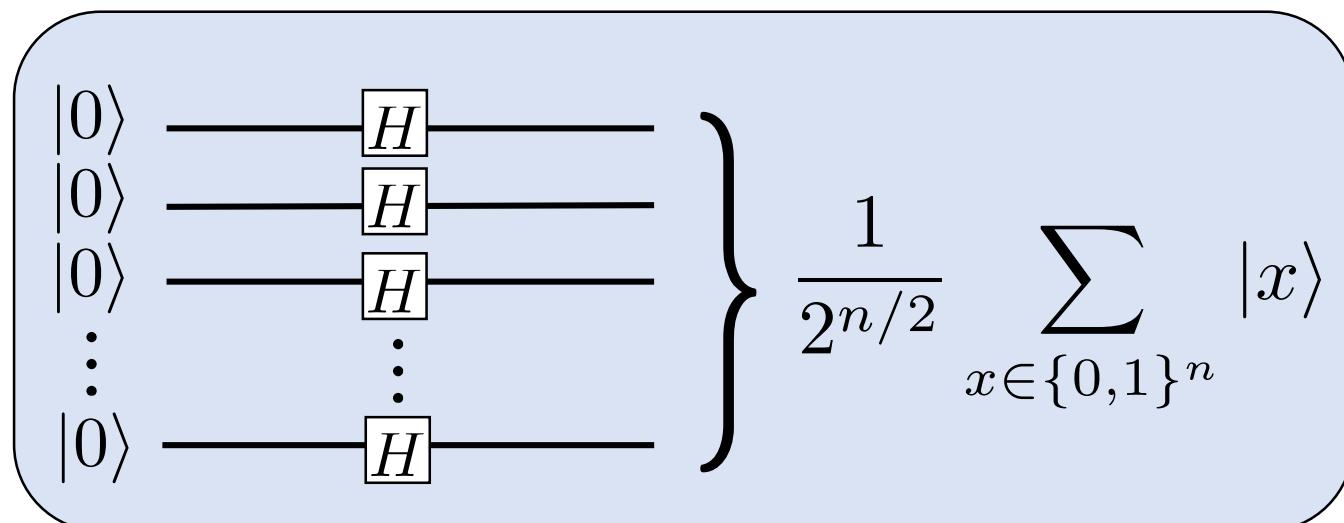


Parallelization

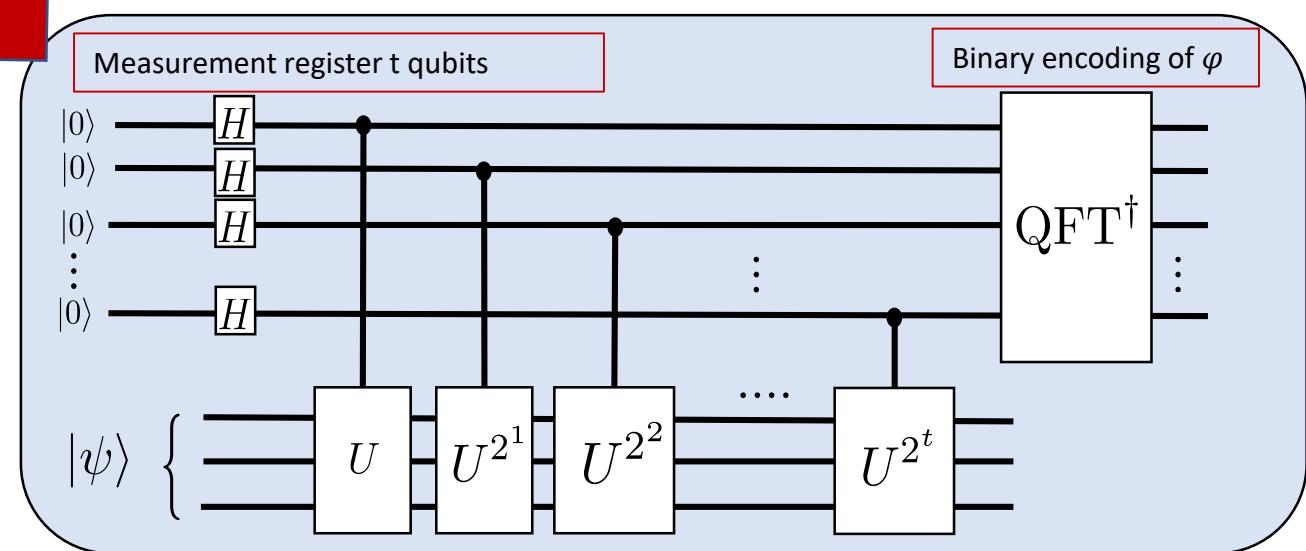
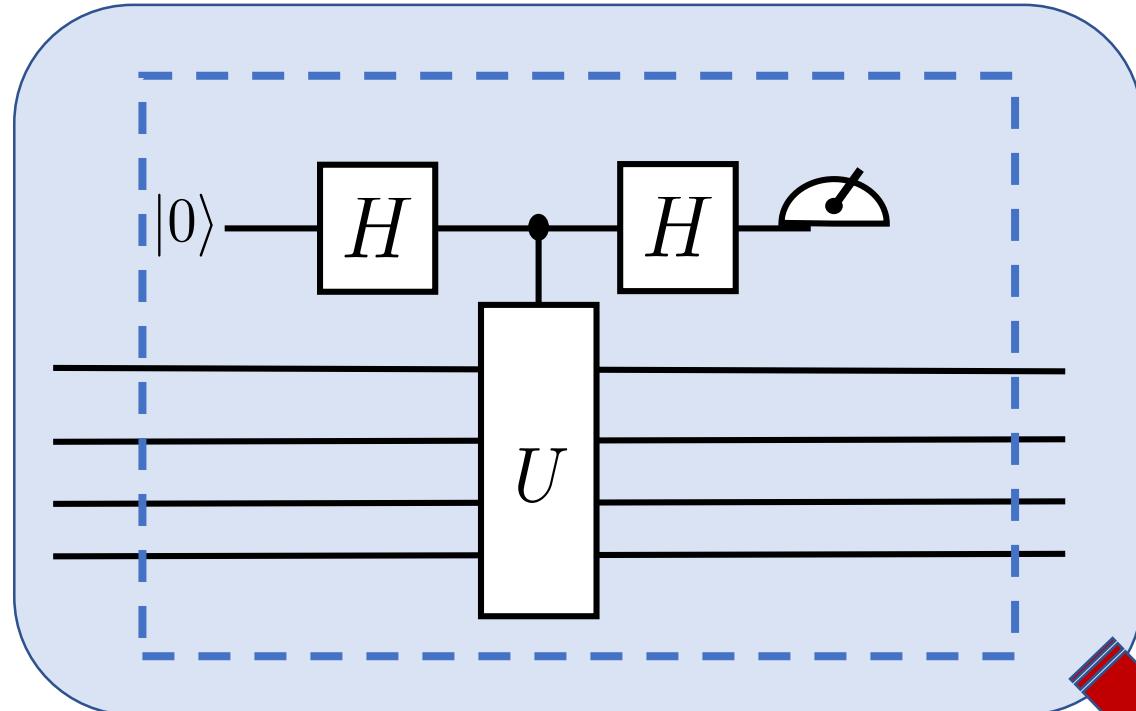
$$|x_1, x_2, \dots, x_{n-1}, x_n\rangle \rightarrow \frac{1}{2^{n/2}}(|0\rangle + e^{2\pi i 0.x_n}|1\rangle) \dots (|0\rangle + e^{2\pi i 0.x_2\dots x_n}|1\rangle)(|0\rangle + e^{2\pi i 0.x_1 x_2 \dots x_n}|1\rangle)$$



$$|x\rangle \xrightarrow{QFT_{2^n}} \frac{1}{2^{n/2}} \sum_{y=0}^{2^n-1} e^{2\pi i xy/2^n} |y\rangle$$



Generalization to N outcomes measurement



Quantum Fourier Transform over \mathbb{Z}_{2^n}

- Quantum circuit: $\Theta(n^2)$ gates.
- Classical FFT: $\Theta(n2^n)$ operations on a 2^n vector.

the first step in phoneme recognition is to Fourier transform the digitized sound. Can we use the quantum Fourier transform to speed up the computation of these Fourier transforms? Unfortunately, the answer is that there is no known way to do this. The problem is that the amplitudes in a quantum computer cannot be directly accessed by measurement. Thus, there is no way of determining the Fourier transformed amplitudes of the original state. Worse still, there is in general no way to efficiently prepare the original state to be Fourier transformed. Thus, finding uses for the quantum Fourier transform is more subtle than we might have hoped. In this and the next chapter we develop several algorithms based upon a more subtle application of the quantum Fourier transform.

Extract from NC
page 220

- A quantum computer does not output the wave-function amplitudes!
 - A quantum computer generates outputs (samples) from a probability distribution $|\psi(x)|^2$
- Quantum mechanics allows you to FFT an exponentially big vector but....
 - encoding an exponentially big classical dataset in a quantum state is highly non-trivial
- Quantum Machine Learning to be practical needs to deal with a similar issue.

References

Reading references

1. Quantum Fourier Transform: NC 5.1,
2. Additional references RdW Ch4, G Ch9

NC ≡ Michael Nielsen and Isaac Chuang, Quantum Computing and Quantum Information
Cambridge University Press (2010)

RdW ≡ Quantum Computing Lecture Notes, Ronald de Wolf, <https://arxiv.org/abs/1907.09415>

G ≡ Introduction to Quantum Computation, Sevag Gharibian, Lectures notes link