

Introduction to Quantum Programming and Semantics

Lecture 10: ZX Calculus

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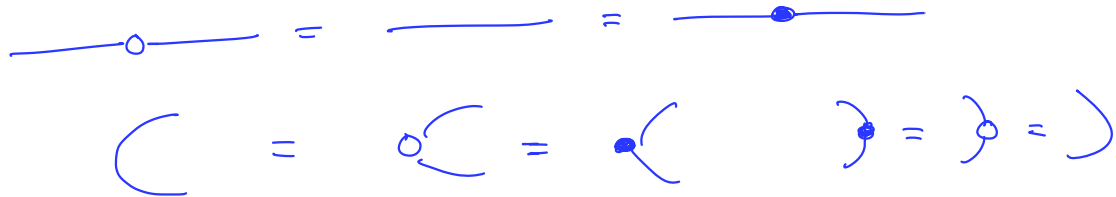
Overview

- ZX Calculus
- Soundness
- Completeness
- Quantum circuits
- Automation

ZX calculus

Rules

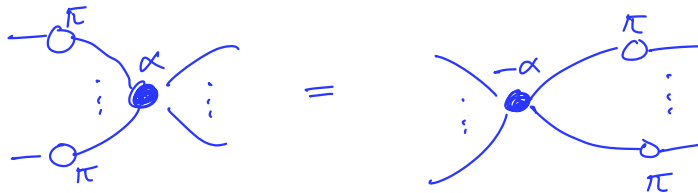
• wire rule:



• spider fusion:

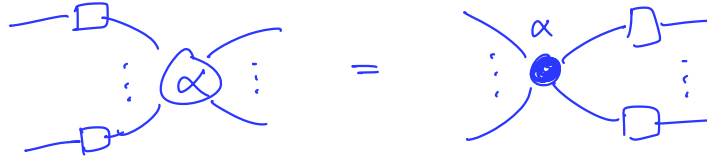


• π -rule:

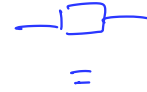


Rules

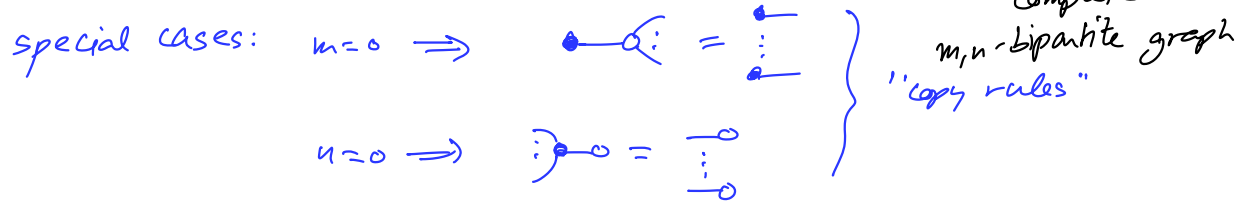
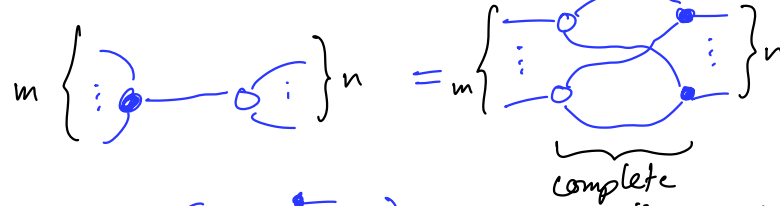
- colour change:



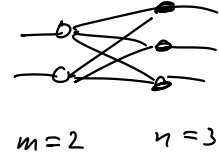
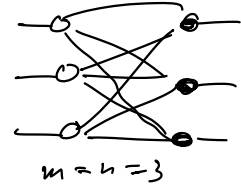
Where



- strong complementarity



e.g.



e.g.:

$$\begin{aligned}
 & \text{CNOT}_1 \text{CNOT}_2 = \pi/2 \pi/2 \pi/2 \pi/2 \pi/2 \pi/2 = \pi/2 \pi/2 \pi \pi/2 \pi/2 = \pi/2 \pi/2 \pi/2 \pi \pi/2 \\
 & = \pi/2 \pi \pi/2 = \text{CNOT} = \text{line}
 \end{aligned}$$

lem:

$$\text{CNOT} = \text{line} \bullet$$

e.g.:

(3 CNOTs)

$$\begin{aligned}
 & \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 \\
 & \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3 = \text{CNOT}_1 \text{CNOT}_2 \text{CNOT}_3
 \end{aligned}$$

(SWAP)

||

Soundness

Semantics

Soundness: if two ZX diagrams d, d' are equal then the matrices $\llbracket d \rrbracket, \llbracket d' \rrbracket$ are equal

$$\llbracket \text{---} \rrbracket = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\llbracket \times \rrbracket = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\llbracket \text{C} \rrbracket = \begin{pmatrix} 1 \\ 0 \\ 0 \\ \vdots \end{pmatrix}$$

$$\llbracket \text{m} \text{---} \text{C} \text{---} \text{n} \rrbracket = \underbrace{|0 \dots 0\rangle}_{\text{n}} \underbrace{\langle 0 \dots 0|}_{\text{m}} + e^{i\alpha} \underbrace{|1 \dots 1\rangle}_{\text{n}} \underbrace{\langle 1 \dots 1|}_{\text{m}} =$$

$$\left(\begin{matrix} \overbrace{1}^{2^m} & & & & \\ & 0 & & 0 & 0 \\ & & 0 & 0 & \\ & 0 & & 0 & \\ & & & & e^{i\alpha} \end{matrix} \right) \Bigg\}^{2^n}$$

$$\llbracket \text{n} \text{---} \text{C}^\alpha \text{---} \text{n} \rrbracket = |+\dots+\rangle \langle +\dots+| + e^{i\alpha} |-\dots-\rangle \langle -\dots-|$$

Completeness

Universality

thm: any n -qubit unitary can be constructed using only
single qubit gates and CNOT gates

cor: any n -qubit unitary can be constructed as a ZX diagram

Approximate universality

then: for any n qubit unitary u and $\varepsilon > 0$,

there exists a ZX diagram d s.t.:

- $\|u - d\| < \varepsilon$
- d only uses phases multiple of $\pi/8$

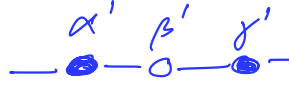
Completeness

thm: if d, d' are ZX diagrams, then:

matrices $\llbracket d \rrbracket = \llbracket d' \rrbracket$ are equal

$\Rightarrow$

$\exists$ legal rewrite from d to d'

needs: $\forall \alpha, \beta, \gamma \exists \alpha', \beta', \gamma'$:  $=$ 

$\hookrightarrow$
in terms of
sin, cos of α, β, γ'

Automation

Quantomatic, PyZX, QuiZX, ZX live

Summary:

- ZX calculus is sound and complete
- ZX calculus is universal
- $\pi/4$ -ZX calculus is approximately universal
- ZX rewriting can be automated