

Introduction to Quantum Programming and Semantics

Lecture 13: Clifford circuits

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Overview

- Clifford circuits and Clifford diagrams
- Graph states with local Clifford operations
- Local complementation, pivoting
- Strong simulation of Clifford circuits
- Synthesis of Clifford circuits

Clifford circuits and Clifford diagrams

Clifford circuits are made from:

$$S = \text{---} \overset{\pi/2}{\circ} \text{---} \\ = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$$

$$H = \text{---} \square \text{---} = \text{---} \overset{\pi/2}{\circ} \overset{\pi/2}{\bullet} \overset{\pi/2}{\circ} \text{---} \\ = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

$$CX = \begin{array}{c} \text{---} \circ \text{---} \\ | \\ \text{---} \bullet \text{---} \end{array} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

e.g.: $Z = \text{---} \overset{\pi}{\circ} \text{---} = \text{---} \overset{\pi/2}{\circ} \overset{\pi/2}{\circ} \text{---} \\ = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$

$$CZ = \begin{array}{c} \text{---} \circ \text{---} \\ | \\ \text{---} \square \text{---} \\ | \\ \text{---} \circ \text{---} \end{array} = \begin{array}{c} \text{---} \circ \text{---} \\ | \\ \text{---} \square \text{---} \bullet \text{---} \square \text{---} \\ | \\ \text{---} \circ \text{---} \end{array}$$

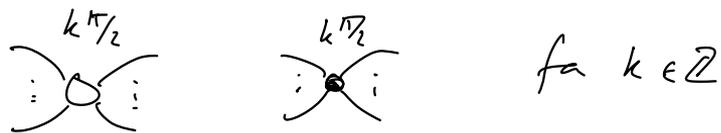
$$CCX = \begin{array}{c} \text{---} \bullet \text{---} \\ | \\ \text{---} \text{---} \\ | \\ \text{---} \oplus \text{---} \end{array}$$

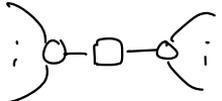
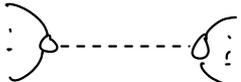
$$X = \text{---} \overset{\pi}{\bullet} \text{---} = \text{---} \square \overset{\pi}{\circ} \square \text{---}$$

$$\sqrt{X} = V = \text{---} \overset{\pi/2}{\bullet} \text{---} = \text{---} \square \overset{\pi/2}{\circ} \square \text{---}$$

but not: $T = \text{---} \overset{\pi/4}{\circ} \text{---}$

Clifford diagrams is a $\mathbb{Z}X$ diagram made of "Clifford spiders"



notation:  \rightsquigarrow  "Hadamard edge"

a $\mathbb{Z}X$ diagram is graph-like if:

- all spiders are \mathbb{Z} spiders
- all edges are Hadamard edges
- no parallel edges or self-loops
- every input/output is connected to a spider

Clifford diagrams

Proposition: every ZX diagram is equal to a graph-like one.

Proof: 1. use colour change rule to eliminate all X spiders

$$\text{---}\bullet\text{---} = \text{---}\square\text{---}\square\text{---}$$

2. use spider fusion to eliminate all non-Hadamard edges

$$\begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \text{---} \begin{array}{c} \circ \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \circ \\ \vdots \end{array}$$

3. eliminate parallel edges

$$\begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \text{---} \begin{array}{c} \circ \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array} \begin{array}{c} \square \\ \vdots \end{array}$$

eliminate self-loops:

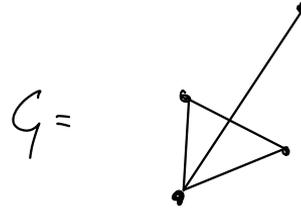
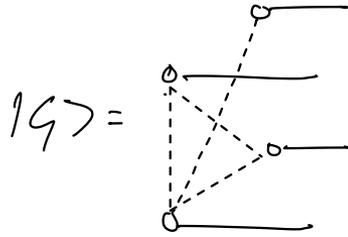
$$\begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} = \begin{array}{c} \vdots \\ \circ \\ \vdots \end{array} \begin{array}{c} \pi/2 \\ \circ \\ \pi/2 \end{array} = \begin{array}{c} \pi/2 \\ \circ \\ \vdots \end{array} = \begin{array}{c} \pi/2 \\ \circ \\ \vdots \end{array} = \begin{array}{c} \pi/2 \\ \circ \\ \vdots \end{array}$$

4. use $\text{---} = \text{---}\circ\text{---}$

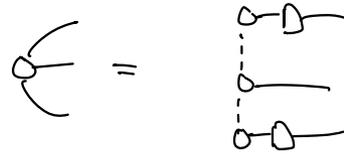
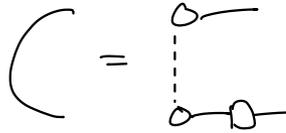
Graph states with local Clifford operations

Graph states are graph-like ZX diagram with no input

e.g.



but not:



def: a graph state with local Cliffords (GSLC)

is a state of form $(U_1 \otimes \dots \otimes U_n) |G\rangle$ for some

graph G and 1-qubit Clifford gates U_i .

Normal form

Def: a Clifford state is a vector of form $C|0 \dots 0\rangle$ for some Clifford circuit C .

Thm: any Clifford state is equal to a GSLC.

Local complementation, pivoting

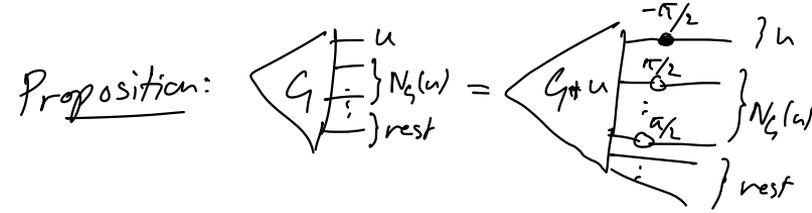
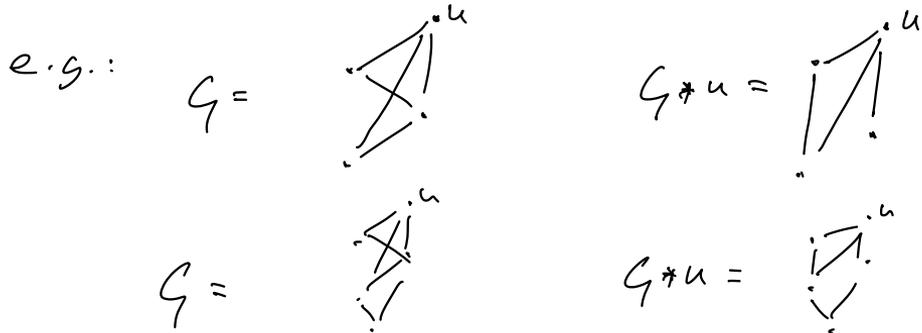
Local complementation of a graph G about vertex u

is a new graph $G \# u = (V, E')$ where

$$\forall v, w \in N_G(u) : \quad (v, w) \in E' \iff (v, w) \notin E$$

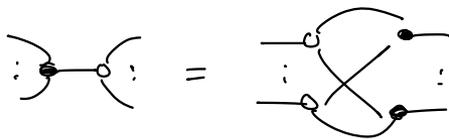
\downarrow
 neighbourhood of u

and $E' \cap (V \setminus N_G(u)) \times V = E \cap (V \setminus N_G(u)) \times V$

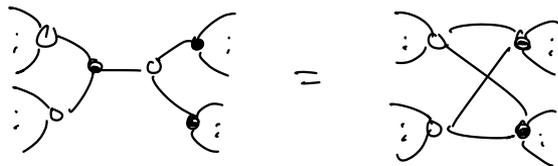


Pivoting

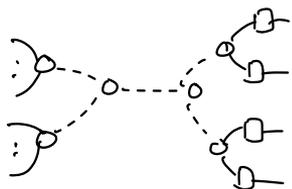
Consider strong complementarity



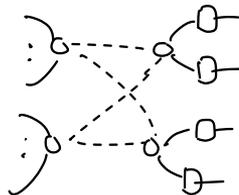
add context



deletes spiders

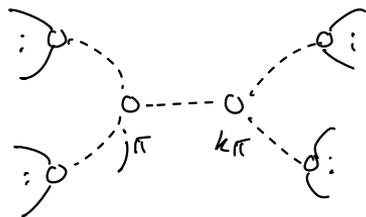


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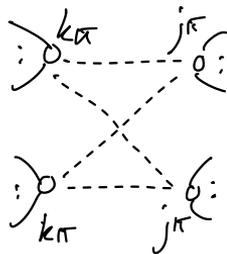


deletes adjacent phase-free spiders

pivot rule



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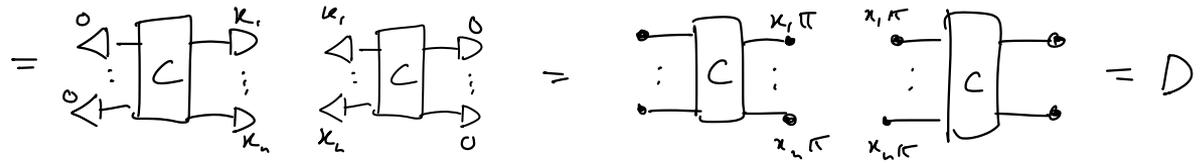


Simulating Clifford circuits

Simulation of Clifford circuits

given Clifford circuit C , compute $\text{Prob}(x_1, \dots, x_n | |\psi\rangle)$, where $|\psi\rangle = C|100\dots 0\rangle$

use Born rule: $\text{Prob}(x_1, \dots, x_n | |\psi\rangle) = |\langle x_1, \dots, x_n | C|10\dots 0\rangle|^2$



algorithm ("Clifford simplification rewrite strat"):

1. convert to graph-like diagram
2. repeat until impossible: local complementation and pivoting
3. remove all isolated $\{0, \pi\}$ -spiders

prop: this terminates

Simulation of Clifford circuits

Prop: algorithm terminates in polynomial time (in $\#$ qubits, and $\#$ gates in C)
 $O((n+k)^3)$ \llcorner_n \llcorner_k

but this is not optimal: doing step 2 optimally $O(n^2 k)$ steps
if $k \gg n$, this makes big difference

idea: 1. avoid big spiders: 

2. apply local complementation and pivoting left-to-right



each step involves at most $O(n)$ spiders (hence $O(n^2)$ wires)

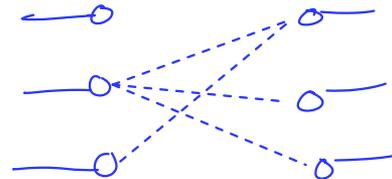
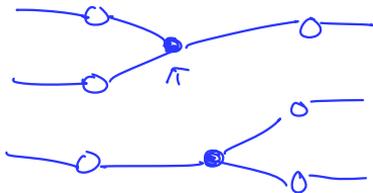
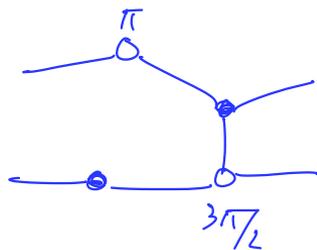
Synthesising Clifford circuits

Synthesis of Clifford circuits

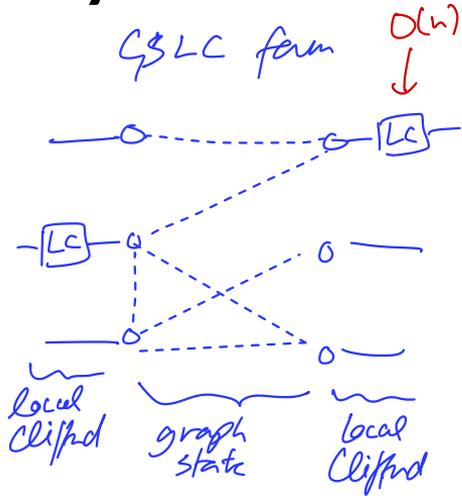
Clifford diagram \longrightarrow "AP-form" \longrightarrow SLC

\hookrightarrow
 all interior spiders
 - have phase $0, \pi$
 - only connected to boundary spiders

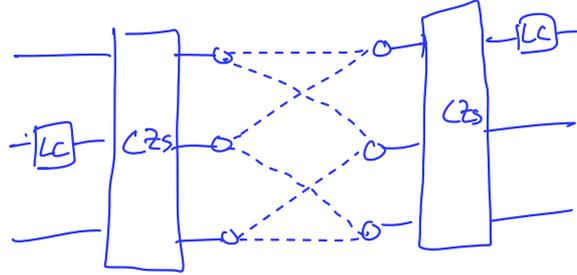
\hookrightarrow
 no internal spiders



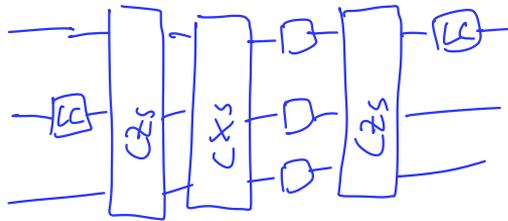
Synthesis of Clifford circuits



unfuse
 \longrightarrow
 CZ gates

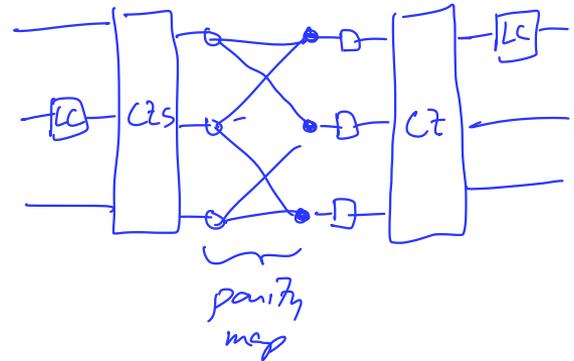


colour change
 \searrow



turn
 parity
 \longleftarrow
 into
 CXs

$O(n^2)$



Any Clifford circuit can be written with at most $O(n^2)$ gates.

Summary:

- ZX diagrams with phases $\pi/2$ are simulable but interesting
- Correspond to Clifford circuits
- Can bring to surprisingly efficient normal form graphically
- Can simulate efficiently graphically