

# Introduction to Quantum Programming and Semantics

## Lecture 3: Composition and tensor product

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# Overview

- OpenQASM

$$\left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_1} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_2} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \right] = \left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_2} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \right] \circ \left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_1} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \right]$$

- Tensor product

$$\left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_1} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \\ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_2} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \end{array} \right] = \left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_1} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \right] \otimes \left[ \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \boxed{C_2} \begin{array}{c} \equiv \\ \equiv \\ \equiv \end{array} \right]$$

- Deutsch-Jozsa: running example

$$\left[ \begin{array}{c} \text{---} \end{array} \boxed{H} \text{---} \right] = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

**OpenQASM**



# OpenQASM

- IBM, but de facto standard  
Braket, Quipper, Tket, many others
- Meant as intermediate representation  
source/target for machines, not humans
- Circuit description  
meant to be executed on classical computer

# Features

- Abstraction:
  - Precise location of qubit doesn't matter (“routing”)
  - Descriptive names for variables
  - Compile-time constants
- Typing
  - qubit, bit, int, angle
  - registers, e.g. qubit[4], int[16]
- Structure
  - conditionals
  - loops
  - subroutines

# Types and gates

- Variables are declared and initialised

```
bit b;  
qubit q;  
reset q;
```

- Single built-in quantum gate **U**

$$U(a, b, c) = \begin{pmatrix} \cos(a/2) & -e^{ic} \sin(a/2) \\ e^{ib} \sin(a/2) & e^{i(b+c)} \cos(a/2) \end{pmatrix}$$

# Subroutines

```
gate X a {  
    U ( pi , 0 , pi ) a;  
}
```

```
gate H a {  
    U ( pi/2 , 0 , pi ) a;  
}
```

# Composition

- Statements separated by semicolons
- Whitespace ignored
- Semantic structure: composition
  - sequential composition of quantum circuits
  - multiplication of unitary matrices
  - sequential composition of string diagrams



# Tensor product

# Registers of multiple qubits

- Size fixed at declaration
- Interpreted as *tensor product*

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} = f : \mathbb{C}^2 \longrightarrow \mathbb{C}^2$$

$$\begin{pmatrix} e & h \\ k & l \end{pmatrix} = g : \mathbb{C}^2 \longrightarrow \mathbb{C}^2$$

$$f \otimes g : \mathbb{C}^2 \otimes \mathbb{C}^2 \longrightarrow \mathbb{C}^2 \otimes \mathbb{C}^2$$

$$\begin{pmatrix} a g & b g \\ c g & d g \end{pmatrix} = \begin{pmatrix} a e & a h & b e & b h \\ a k & a l & b k & b l \\ c e & c h & d e & d h \\ c k & c l & d k & d l \end{pmatrix}$$

*"Kronecker product"*

Single qubit:  $|\psi\rangle \in \mathbb{C}^2$

$n$  qubits:  $\underbrace{(\mathbb{C}^2) \otimes \dots \otimes (\mathbb{C}^2)}_{n \text{ times}}$

$0$  qubits:  $(\mathbb{C}^n) = \mathbb{C}^1 = \mathbb{C}$

$$\begin{array}{ccc}
 (v, w) & \xrightarrow{\quad} & v \otimes w \\
 V \times W & \xrightarrow{\text{bilinear}} & V \otimes W = \left\{ \sum_{i=1}^n z_i v_i \otimes w_i \mid \begin{array}{l} z_i \in \mathbb{C} \\ v_i \in V \\ w_i \in W \end{array} \right\} \\
 & \searrow \text{bilinear} & \downarrow \exists! \text{ linear} \\
 & & X
 \end{array}$$

"Tensor products make bilinear maps linear"

# Tensor product interacts with composition

- Interchange law

$$\llbracket \text{---} f_r \text{---} f_l \text{---} \rrbracket \otimes \llbracket \text{---} g_r \text{---} g_l \text{---} \rrbracket$$

$$= \left[ \begin{array}{c} \text{---} f_r \text{---} \\ \text{---} g_r \text{---} \end{array} \right] ; \left[ \begin{array}{c} \text{---} f_l \text{---} \\ \text{---} g_l \text{---} \end{array} \right]$$

$$= \left[ \begin{array}{c} \text{---} f_r \text{---} f_l \text{---} \\ \text{---} g_r \text{---} g_l \text{---} \end{array} \right]$$

sequential composition is

associative

$$h \circ (g \circ f) = (h \circ g) \circ f$$

"after"

not commutative

$$g \circ f \neq f \circ g$$

parallel composition

associative

$$h \otimes (g \otimes f) = (h \otimes g) \otimes f$$

symmetry

$$V \otimes W \simeq W \otimes V$$

# Controlled gates

- Modifier

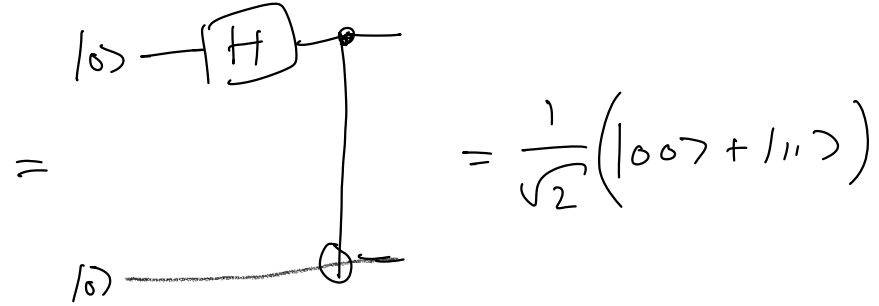
`ctrl @ U (a, b, c) q_reg [0], q_reg [1];`

- Interpreted semantically as matrix

$$1 \oplus U(a, b, c) = \begin{array}{cc|cc} & |00\rangle & |01\rangle & |10\rangle & |11\rangle \\ \begin{array}{c} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{array} & \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ \hline 0 & 0 & - & - \\ 0 & 0 & - & - \end{pmatrix} \end{array}$$

# Example

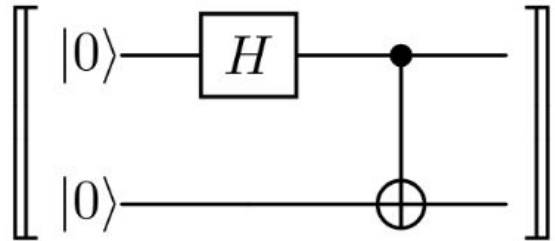
```
gate H a {  
    U (pi/2, 0, pi) a;  
}  
  
gate CX a, b {  
    ctrl @ U (pi, 0, pi) a, b;  
}  
  
qubit [2] = q_reg;  
reset q_reg;  
H q_reg [0];  
CX q_reg [0], q_reg [1];
```



# Example

```
gate H a {  
    U (pi/2, 0, pi) a;  
}  
  
gate CX a, b {  
    ctrl @ U (pi, 0, pi) a, b;  
}
```

```
qubit [2] = q_reg;  
reset q_reg;  
H q_reg [0];  
CX q_reg [0], q_reg [1];
```


$$\left[ \begin{array}{c} |0\rangle \text{---} \boxed{H} \text{---} \bullet \\ |0\rangle \text{---} \oplus \end{array} \right] = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$



**Deutsch-Jozsa**

# All-or-nothing oracular promise problems

- Decide on a solution without relying on approximation
- Input is provided as oracle
- Relies on promise about 'global behaviour' of input
- Quantum algorithm faster than best-known classical algorithm
- E.g. Shor, Grover, hidden subgroup

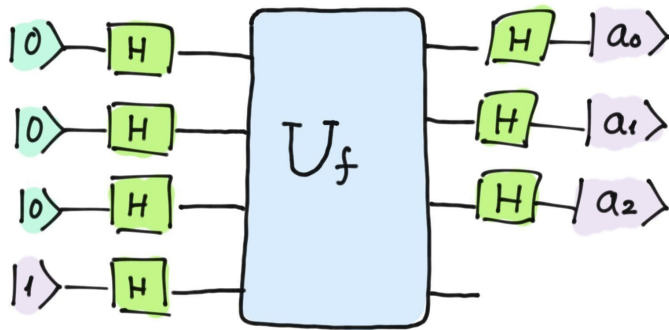
# Setting

- Input: 2-colouring  $f: \{0,1\}^n \rightarrow \{0,1\}$  of bitstrings
- Promise:  $f$  is either
  - Constant: any input gets same colour
  - Balanced: 0 on half the inputs, 1 on the other half
- Task: find out which is the case

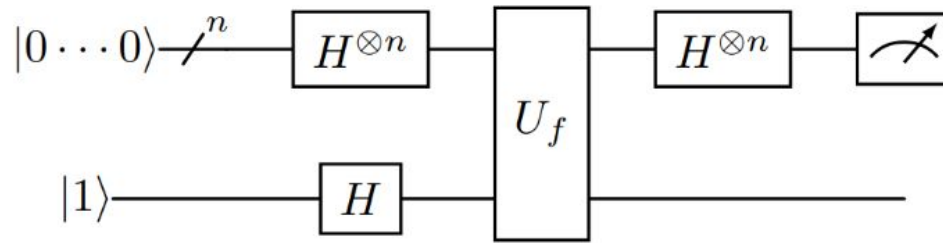
# Oracles

- Can ask  $f$  for its value on some bitstring
- *Bennett's trick* turns  $f$  into unitary gate

$$U_f |xy\rangle = |x\rangle \otimes |y \text{ XOR } f(x)\rangle$$



# Deutsch-Jozsa circuit

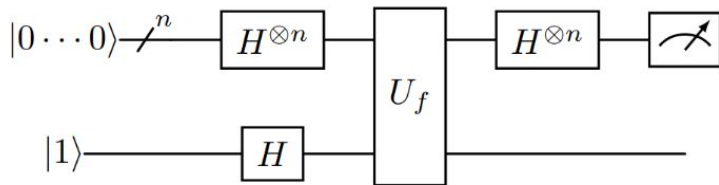


# Example oracle

```
gate Oracle x y z {  
    CX x z;  
    CX y z;  
}
```

$xyz$	$x \text{ XOR } y$	$z \text{ XOR } (x \text{ XOR } y)$
000	0	0
001	0	1
010	1	1
011	1	0
100	1	1
101	1	0
110	0	0
111	0	1

# Implementation



```
// declare three qubits
qreg x;
qreg y;
qreg z;
// set the three qubits to  $|0\rangle$ ,  $|0\rangle$ , and  $|1\rangle$ 
reset x;
reset y;
reset z;
X z;
// apply Hadamard to all three qubits
H x;
H y;
H z;
// apply the oracle
Oracle x y z;
// apply Hadamard to the first two qubits
H x;
H y;
// measure the first two qubits (will discuss later)
bit a = measure x;
bit b = measure y;
```

# Summary:

- OpenQASM low-level de-facto standard circuit description language
- To interpret it, semantics needs both sequential and parallel composition
- Tensor products of matrices
- Deutsch-Jozsa prototypically solves all-or-nothing oracular promise problem