

Isabelle/HOL Exercises

Logic and Sets

Propositional Logic

In this exercise, we will prove some lemmas of propositional logic with the aid of a calculus of natural deduction.

For the proofs, you may only use

- the following lemmas:

notI: $(A \Longrightarrow \text{False}) \Longrightarrow \neg A$,

notE: $[\neg A; A] \Longrightarrow B$,

conjI: $[[A; B] \Longrightarrow A \wedge B$,

conjE: $[A \wedge B; [[A; B] \Longrightarrow C]] \Longrightarrow C$,

disjI1: $A \Longrightarrow A \vee B$,

disjI2: $A \Longrightarrow B \vee A$,

disjE: $[A \vee B; A \Longrightarrow C; B \Longrightarrow C] \Longrightarrow C$,

impI: $(A \Longrightarrow B) \Longrightarrow A \longrightarrow B$,

impE: $[A \longrightarrow B; A; B \Longrightarrow C] \Longrightarrow C$,

mp: $[A \longrightarrow B; A] \Longrightarrow B$

iffI: $[A \Longrightarrow B; B \Longrightarrow A] \Longrightarrow A = B$,

iffE: $[A = B; [A \longrightarrow B; B \longrightarrow A] \Longrightarrow C] \Longrightarrow C$

classical: $(\neg A \Longrightarrow A) \Longrightarrow A$

- the proof methods *rule*, *erule* and *assumption*.

Prove:

lemma *I*: " $A \longrightarrow A$ "

lemma " $A \wedge B \longrightarrow B \wedge A$ "

lemma " $(A \wedge B) \longrightarrow (A \vee B)$ "

lemma " $((A \vee B) \vee C) \longrightarrow A \vee (B \vee C)$ "

lemma *K*: " $A \longrightarrow B \longrightarrow A$ "

lemma " $(A \vee A) = (A \wedge A)$ "

lemma *S*: " $(A \longrightarrow B \longrightarrow C) \longrightarrow (A \longrightarrow B) \longrightarrow A \longrightarrow C$ "

lemma " $(A \longrightarrow B) \longrightarrow (B \longrightarrow C) \longrightarrow A \longrightarrow C$ "

lemma " $\neg \neg A \longrightarrow A$ "

lemma " $A \longrightarrow \neg \neg A$ "

lemma " $(\neg A \longrightarrow B) \longrightarrow (\neg B \longrightarrow A)$ "

lemma " $(A \longrightarrow B) \longrightarrow A \longrightarrow A$ "

lemma " $A \vee \neg A$ "

lemma " $\neg (A \wedge B) = (\neg A \vee \neg B)$ "